

The Density of Positive Entries of a Linear Recurrence

Edon Kelmendi

Max Planck Institute for Software Systems
Saarbrücken, Germany

I. The Problem

II. The Theorems

**III. The Example, or First
Observation**

IV. The Proof

V. The Open Problem

I. The Problem

$x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$

while true do

$x \leftarrow 4x + 3y$

$y \leftarrow 4y - 3x$

$z \leftarrow 5z$

if $y + z > 0$ then

Region A

else

Region B

end if

end while

$x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$

while true do

$x \leftarrow 4x + 3y$

$y \leftarrow 4y - 3x$

$z \leftarrow 5z$

if $y + z > 0$ **then**

Region A

else

Region B

end if

end while

• arbitrary number of variables ranging over integers

• linear updates

• polynomial inequality

$x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$

while true do

▶ $x \leftarrow 4x + 3y$

$y \leftarrow 4y - 3x$

$z \leftarrow 5z$

if $y + z > 0$ then

Region A

else

Region B

end if

end while

$x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$

while true do

$x \leftarrow 4x + 3y$

▶ $y \leftarrow 4y - 3x$

$z \leftarrow 5z$

if $y + z > 0$ then

Region A

else

Region B

end if

end while

$x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$

while true do

$x \leftarrow 4x + 3y$

$y \leftarrow 4y - 3x$

▶ $z \leftarrow 5z$

if $y + z > 0$ then

Region A

else

Region B

end if

end while

$x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$

while true do

$x \leftarrow 4x + 3y$

$y \leftarrow 4y - 3x$

$z \leftarrow 5z$

▶ **if $y + z > 0$ then**

Region A

else

Region B

end if

end while

$x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$

while true do

$x \leftarrow 4x + 3y$

$y \leftarrow 4y - 3x$

$z \leftarrow 5z$

if $y + z > 0$ then



► **Region A**

else

Region B

end if

end while



$x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$

while true do

$x \leftarrow 4x + 3y$

$y \leftarrow 4y - 3x$

$z \leftarrow 5z$

if $y + z > 0$ then



► **Region A**

else

Region B

end if

end while

$x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$

while true do

$x \leftarrow 4x + 3y$

$y \leftarrow 4y - 3x$

$z \leftarrow 5z$

if $y + z > 0$ then



▶ **Region A**

else

Region B

end if

end while



$x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$

while true do

$x \leftarrow 4x + 3y$

$y \leftarrow 4y - 3x$

$z \leftarrow 5z$

if $y + z > 0$ then

Region A

else

► Region B

end if

end while



$x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$

while true do

$x \leftarrow 4x + 3y$

$y \leftarrow 4y - 3x$

$z \leftarrow 5z$

if $y + z > 0$ then

Region A

else

► Region B

end if

end while



$x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$

while true do

$x \leftarrow 4x + 3y$

$y \leftarrow 4y - 3x$

$z \leftarrow 5z$

if $y + z > 0$ then

Region A

else

Region B

end if



end while



```
 $x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$   
while true do  
   $x \leftarrow 4x + 3y$   
   $y \leftarrow 4y - 3x$   
   $z \leftarrow 5z$   
  if  $y + z > 0$  then  
    Region A  
  else  
    Region B  
  end if  
end while
```



Decision questions:

1. Is Region A reached?
(Is there at least one  ?)
2. Is Region A reached infinitely often?
(Are there infinitely many  ?)


```

 $x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$ 
while true do
   $x \leftarrow 4x + 3y$ 
   $y \leftarrow 4y - 3x$ 
   $z \leftarrow 5z$ 
  if  $y + z > 0$  then
    Region A
  else
    Region B
  end if
end while

```



Decision questions:

1. Is **Region A** reached?
(Is there at least one ■ ?)
 - Known as the **positivity problem**;
at least as hard as Skolem's problem
2. Is Region A reached infinitely often?
(Are there infinitely many ■ ?)
 - Known as the **ultimate positivity problem**;
also open & difficult

```
x ← 0; y ← 6; z ← 4
while true do
  x ← 4x + 3y
  y ← 4y - 3x
  z ← 5z
  if y + z > 0 then
    Region A
  else
    Region B
  end if
end while
```



Decision questions:

1. Is **Region A** reached?
(Is there at least one ■ ?)
 - Known as the **positivity problem**;
at least as hard as Skolem's problem
2. Is Region A reached infinitely often?
(Are there infinitely many ■ ?)
 - Known as the **ultimate positivity problem**;
also open & difficult

In this paper:

3. How much more frequent are ■ compared to ■?

$x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$

while true do

$x \leftarrow 4x + 3y$

$y \leftarrow 4y - 3x$

$z \leftarrow 5z$

if $y + z > 0$ **then**

Region A

else

Region B

end if

end while



Decision questions:

1. Is Region A reached?

(Is there at least one ■ ?)

- Known as the **positivity problem**;
at least as hard as Skolem's problem

2. Is Region A reached infinitely often?

(Are there infinitely many ■ ?)

- Known as the **ultimate positivity problem**;
also open & difficult

In this paper:

3. How much more frequent are ■ compared to ■?

Set of ■

1. Is it empty?
2. Is it infinite?
3. How big is it inside \mathbb{N} ?

Density

$$S \subset \mathbb{N} \quad \mathcal{D}_\circ(S) := \liminf_{n \rightarrow \infty} \frac{|\{1, 2, \dots, n\} \cap S|}{n}$$

What proportion of first n numbers are in S ($n \rightarrow \infty$) ?

Examples:

- $\mathcal{D}_\circ(\text{Finite Set}) = 0$, $\mathcal{D}_\circ(\text{co-Finite Set}) = 1$
- $\mathcal{D}_\circ(\text{Arithmetic progression of length } k) = 1/k$,
- $\mathcal{D}_\circ(\text{Geometric progression}) = 0$,
- $\mathcal{D}_\circ(\text{Primes}) = 0$, (Prime Number Theorem).

Density

$$S \subset \mathbb{N} \quad \mathcal{D}_\circ(S) := \liminf_{n \rightarrow \infty} \frac{|\{1, 2, \dots, n\} \cap S|}{n}$$

What proportion of first n numbers are in S ($n \rightarrow \infty$)?

$\mathcal{D}_\circ(X) = 0 \Rightarrow X$ is sparse

$\mathcal{D}_\circ(X) = 1 \Rightarrow X$ is very dense

if $\mathcal{D}_\circ(X) = \mathcal{D}^\circ(X)$, denote by $\mathcal{D}(X)$

$x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$

while true do

$x \leftarrow 4x + 3y$

$y \leftarrow 4y - 3x$

$z \leftarrow 5z$

if $y + z > 0$ then

Region A


else

Region B

end if

end while



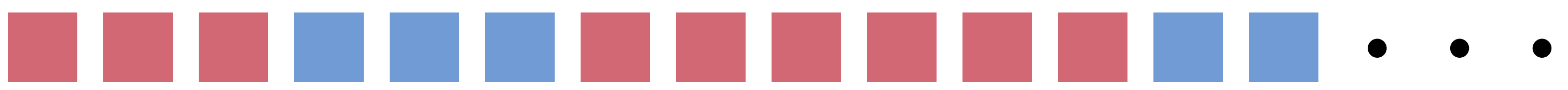
Number of  in first n entries

n , $n \in \mathbb{N}$

```

 $x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$ 
while true do
   $x \leftarrow 4x + 3y$ 
   $y \leftarrow 4y - 3x$ 
   $z \leftarrow 5z$ 
  if  $y + z > 0$  then
    Region A
  else
    Region B
  end if
end while

```



$$\lim_{n \rightarrow \infty} \frac{\text{Number of } \blacksquare \text{ in first } n \text{ entries}}{n}$$

Exists due to:

Jason P Bell and Stefan Gerhold. On the positivity set of a linear recurrence sequence. *Israel Journal of Mathematics*, 157(1):333–345, 2007.

Denote it by \mathcal{D} .
 Main character of the story

II. The Theorems

Theorems

Theorem 1. “ $\mathcal{D} = 0?$ ” is decidable.

(so is “ $\mathcal{D} = 1?$ ” by symmetry)

Theorem 1a. For diagonalisable update matrices:

$\mathcal{D} = 0 \Leftrightarrow$ finitely many ■

Theorem 2. \mathcal{D} can be computed to arbitrary precision.

Theorem 3. “ $\mathcal{D} \in \mathbb{Q}?$ ” is decidable,

when there are at most three dominant eigenvalues.

Theorems

Theorem 1. “ $\mathcal{D} = 0?$ ” is decidable.

(so is “ $\mathcal{D} = 1?$ ” by symmetry)

Theorem 1a. For diagonalisable update matrices:

$\mathcal{D} = 0 \Leftrightarrow$ finitely many ■

Theorem 2. \mathcal{D} can be computed to arbitrary precision.

Theorem 3. “ $\mathcal{D} \in \mathbb{Q}?$ ” is decidable,

when there are at most three dominant eigenvalues.

Theorems

Theorem 1. “ $\mathcal{D} = 0?$ ” is decidable.

(so is “ $\mathcal{D} = 1?$ ” by symmetry)

Theorem 1a. For diagonalisable update matrices:

$\mathcal{D} = 0 \Leftrightarrow$ finitely many ■

Theorem 2. \mathcal{D} can be computed to arbitrary precision.

Theorem 3. “ $\mathcal{D} \in \mathbb{Q}?$ ” is decidable,

when there are at most three dominant eigenvalues.

Theorems

Theorem 1. “ $\mathcal{D} = 0?$ ” is decidable.

(so is “ $\mathcal{D} = 1?$ ” by symmetry)

Theorem 1a. For diagonalisable update matrices:

$\mathcal{D} = 0 \Leftrightarrow$ finitely many ■

Theorem 2. \mathcal{D} can be computed to arbitrary precision.

Theorem 3. “ $\mathcal{D} \in \mathbb{Q}?$ ” is decidable,

when there are at most three dominant eigenvalues.

```
x ← 0; y ← 6; z ← 4
while true do
  x ← 4x + 3y
  y ← 4y - 3x
  z ← 5z
  if y + z > 0 then
    Region A
  else
    Region B
  end if
end while
```



To Summarise:

- We can decide if **Region A** is entered not too rarely
- We can say a lot about the asymptotic frequency of entering **Region A/Region B**

III. The Example, or First Observation

```
 $x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$   
while true do  
   $x \leftarrow 4x + 3y$   
   $y \leftarrow 4y - 3x$   
   $z \leftarrow 5z$   
  if  $y + z > 0$  then  
    Region A  
  else  
    Region B  
  end if  
end while
```

How frequently is Region A entered?

```
 $x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$   
while true do  
   $x \leftarrow 4x + 3y$   
   $y \leftarrow 4y - 3x$   
   $z \leftarrow 5z$   
  if  $y + z > 0$  then  
    Region A  
  else  
    Region B  
  end if  
end while
```

How frequently is Region A entered?

$$0.732279 \dots = \frac{\cos^{-1}(-2/3)}{\pi}$$

$x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$

while true do

$x \leftarrow 4x + 3y$
 $y \leftarrow 4y - 3x$
 $z \leftarrow 5z$

if $y + z > 0$ **then**

Region A

else

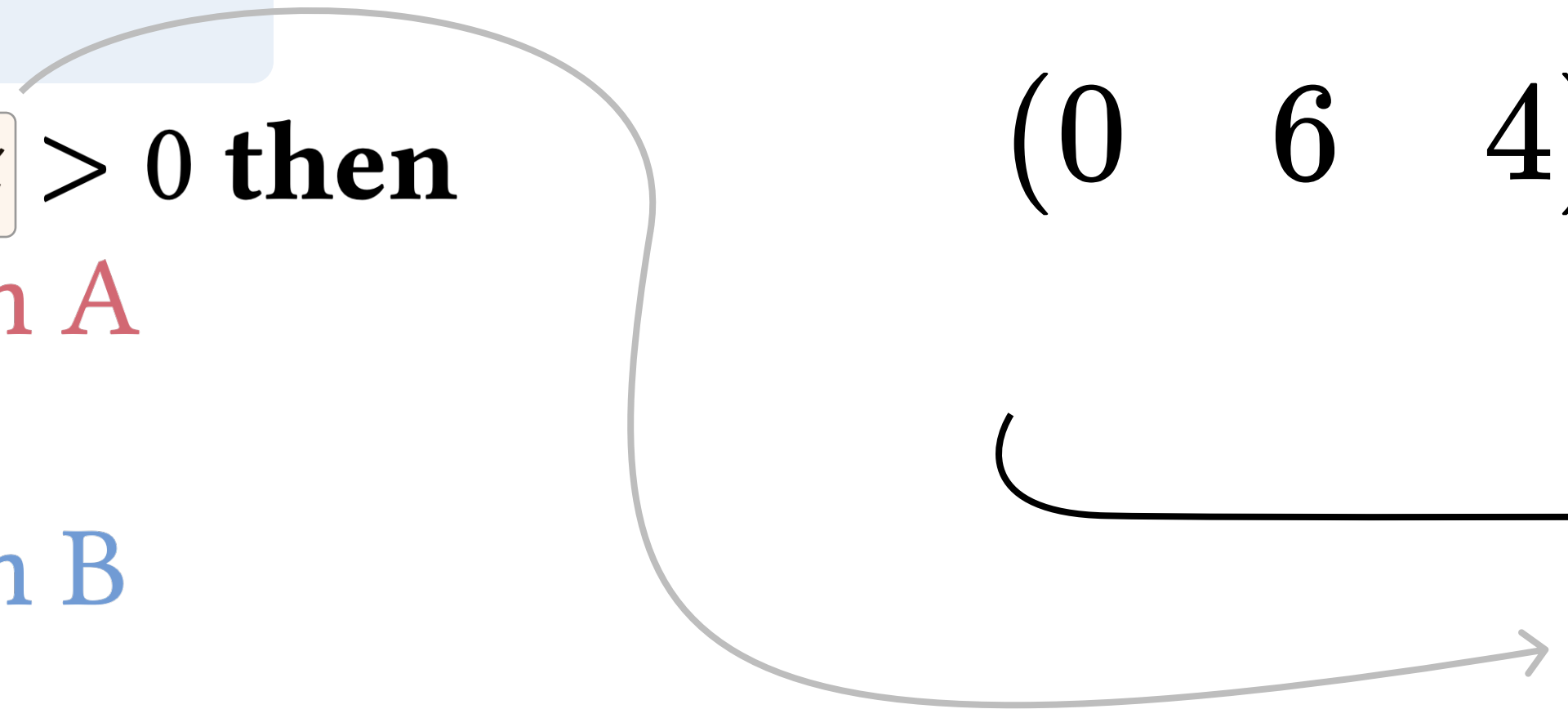
Region B

end if

end while

$$(0 \quad 6 \quad 4) \cdot \begin{pmatrix} 4 & -3 & 0 \\ 3 & 4 & 0 \\ 0 & 0 & 5 \end{pmatrix}^n \cdot \begin{pmatrix} 0 \\ 1 \\ 1 \end{pmatrix}$$

$y+z$ in the n th loop iteration



```

x ← 0; y ← 6; z ← 4
while true do
  x ← 4x + 3y
  y ← 4y - 3x
  z ← 5z
  if y + z > 0 then
    Region A
  else
    Region B
  end if
end while

```

$$(0 \ 6 \ 4) \cdot \begin{pmatrix} 4 & -3 & 0 \\ 3 & 4 & 0 \\ 0 & 0 & 5 \end{pmatrix}^n \cdot \begin{pmatrix} 0 \\ 1 \\ 1 \end{pmatrix} > 0$$

y+z in the nth loop iteration

$$5^n \begin{pmatrix} 4/5 & -3/5 & 0 \\ 3/5 & 4/5 & 0 \\ 0 & 0 & 1 \end{pmatrix}^n$$

2D rotation

$$5^n \begin{pmatrix} 4/5 & -3/5 & 0 \\ 3/5 & 4/5 & 0 \\ 0 & 0 & 1 \end{pmatrix}^n$$

2D rotation

$$\begin{aligned} r(\cos \theta, \sin \theta) \cdot \begin{pmatrix} \cos \varphi & \sin \varphi \\ -\sin \varphi & \cos \varphi \end{pmatrix} &= r(\cos \theta \cos \varphi - \sin \theta \sin \varphi, \cos \theta \sin \varphi + \sin \theta \cos \varphi) \\ &= r(\cos(\theta + \varphi), \sin(\theta + \varphi)) \end{aligned}$$

```

x ← 0; y ← 6; z ← 4
while true do
  x ← 4x + 3y
  y ← 4y - 3x
  z ← 5z
  if y + z > 0 then
    Region A
  else
    Region B
  end if
end while

```

$$(0 \ 6 \ 4) \cdot \begin{pmatrix} 4 & -3 & 0 \\ 3 & 4 & 0 \\ 0 & 0 & 5 \end{pmatrix}^n \cdot \begin{pmatrix} 0 \\ 1 \\ 1 \end{pmatrix} > 0$$

y+z in the nth loop iteration

$$5^n \begin{pmatrix} 4/5 & -3/5 & 0 \\ 3/5 & 4/5 & 0 \\ 0 & 0 & 1 \end{pmatrix}^n = 5^n \begin{pmatrix} \cos n\varphi & -\sin n\varphi & 0 \\ \sin n\varphi & \cos n\varphi & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

Rotation in the first two coordinates by $\varphi = \cos^{-1} 4/5$

Rotation in the first two coordinates by $\varphi = \cos^{-1} 4/5$

```
x ← 0; y ← 6; z ← 4
```

```
while true do
```

```
  x ← 4x + 3y
```

```
  y ← 4y - 3x
```

```
  z ← 5z
```

```
  if y + z > 0 then
```

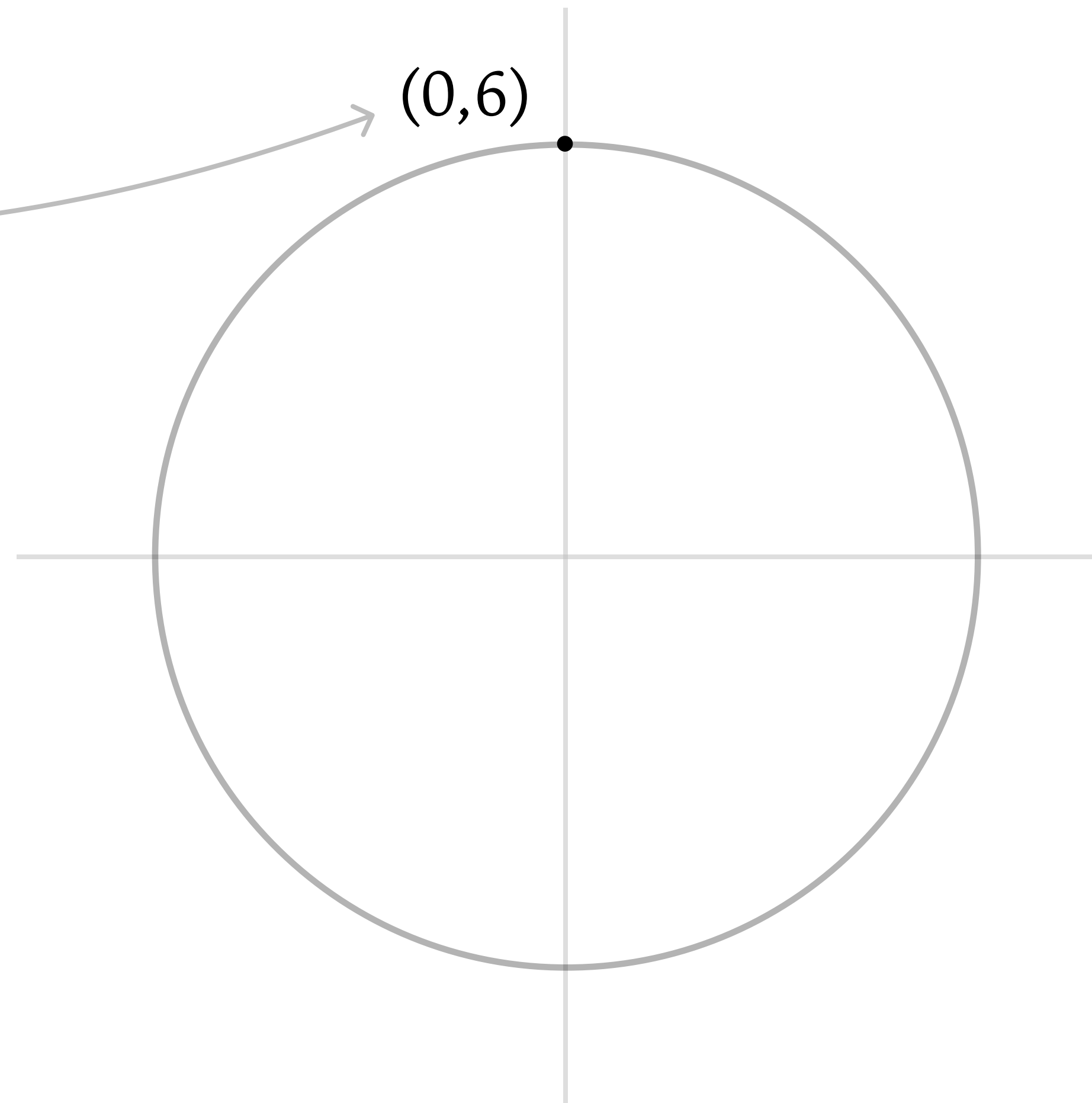
```
    Region A
```

```
  else
```

```
    Region B
```

```
  end if
```

```
end while
```



Rotation in the first two coordinates by $\varphi = \cos^{-1} 4/5$

```
x ← 0; y ← 6; z ← 4
```

```
while true do
```

```
  x ← 4x + 3y
```

```
  y ← 4y - 3x
```

```
  z ← 5z
```

```
  if y + z > 0 then
```

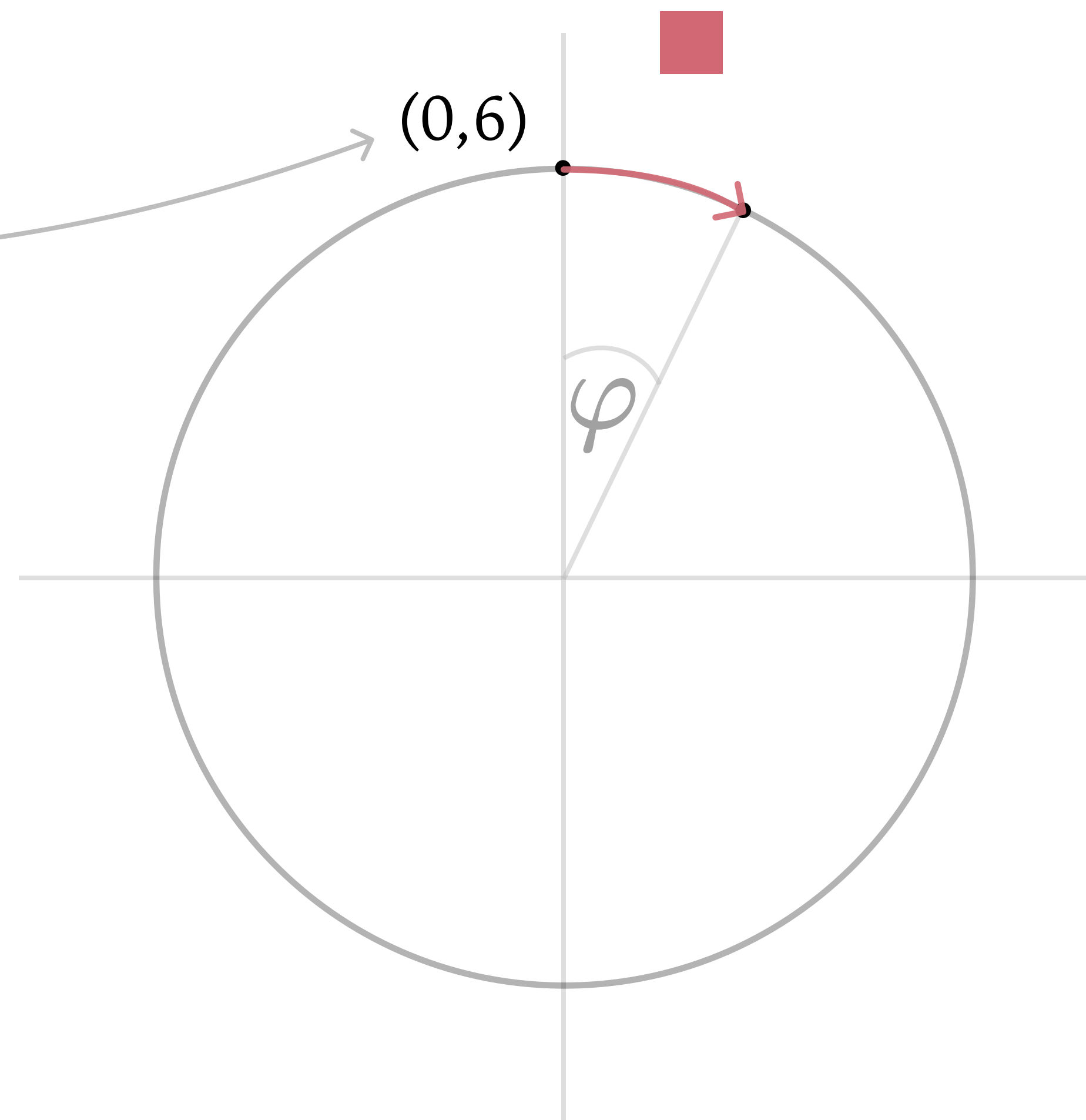
```
    Region A
```

```
  else
```

```
    Region B
```

```
  end if
```

```
end while
```



Rotation in the first two coordinates by $\varphi = \cos^{-1} 4/5$

```
x ← 0; y ← 6; z ← 4
```

```
while true do
```

```
  x ← 4x + 3y
```

```
  y ← 4y - 3x
```

```
  z ← 5z
```

```
  if y + z > 0 then
```

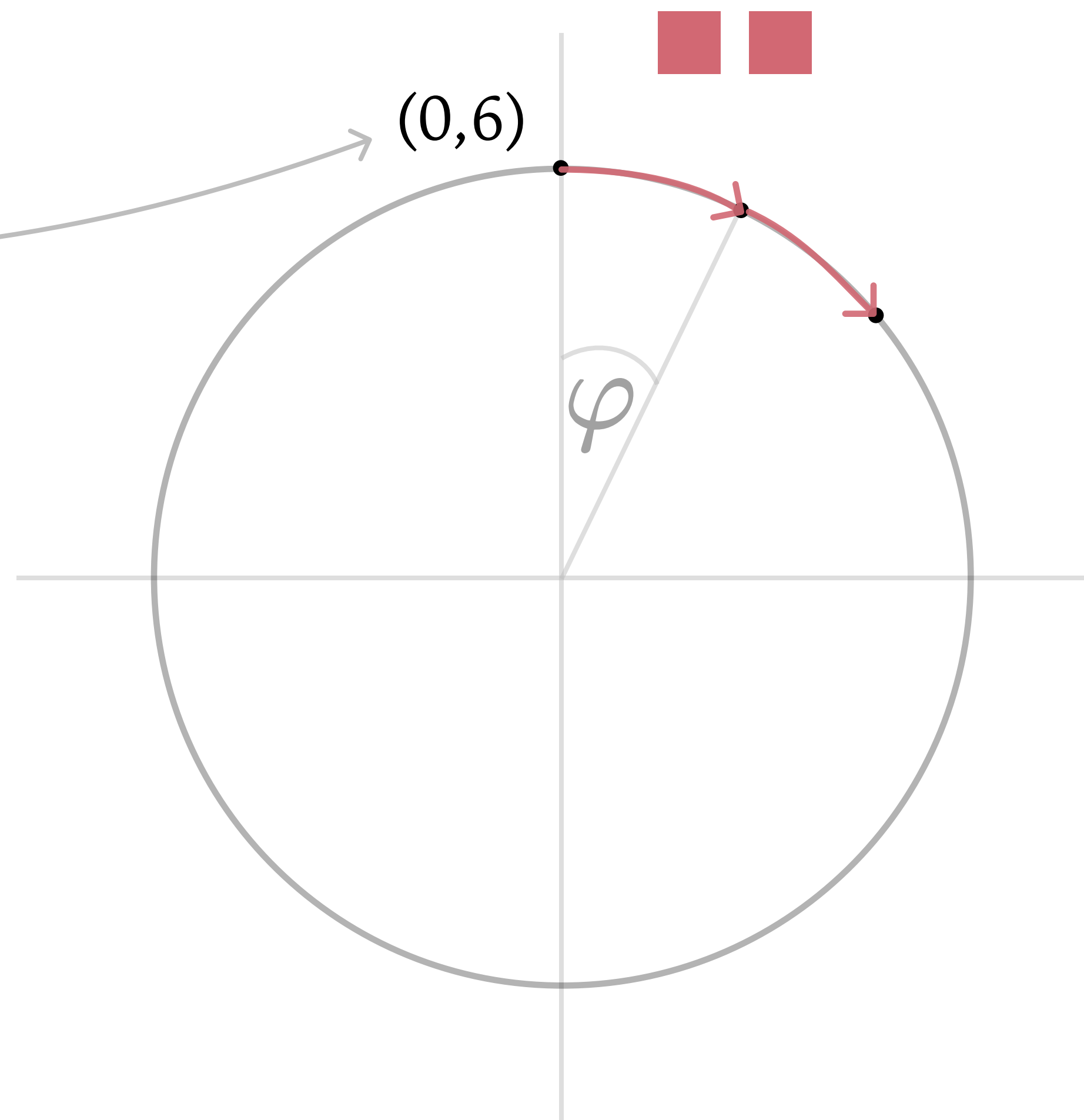
```
    Region A
```

```
  else
```

```
    Region B
```

```
  end if
```

```
end while
```



Rotation in the first two coordinates by $\varphi = \cos^{-1} 4/5$

```
x ← 0; y ← 6; z ← 4
```

```
while true do
```

```
  x ← 4x + 3y
```

```
  y ← 4y - 3x
```

```
  z ← 5z
```

```
  if y + z > 0 then
```

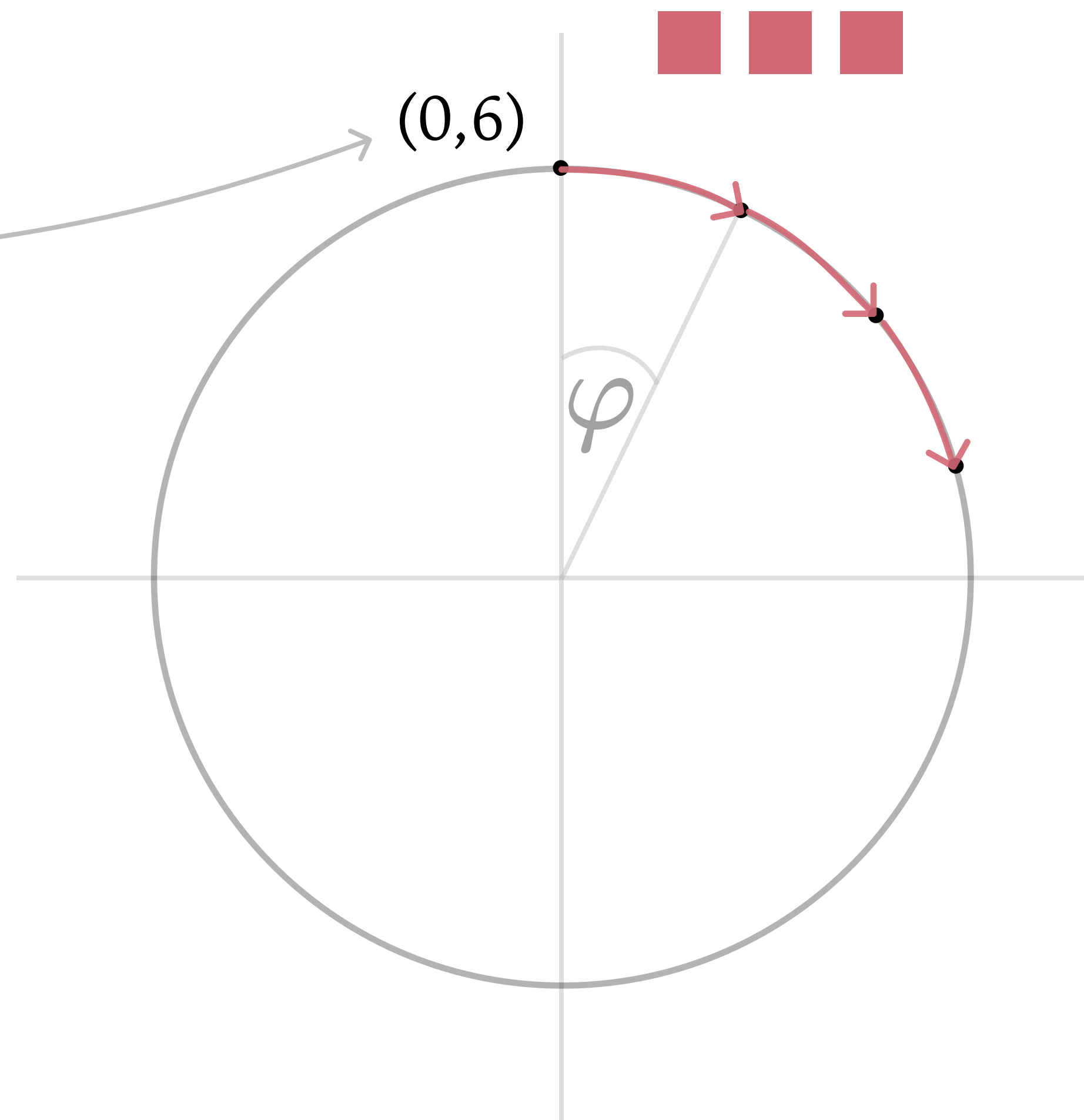
```
    Region A
```

```
  else
```

```
    Region B
```

```
  end if
```

```
end while
```

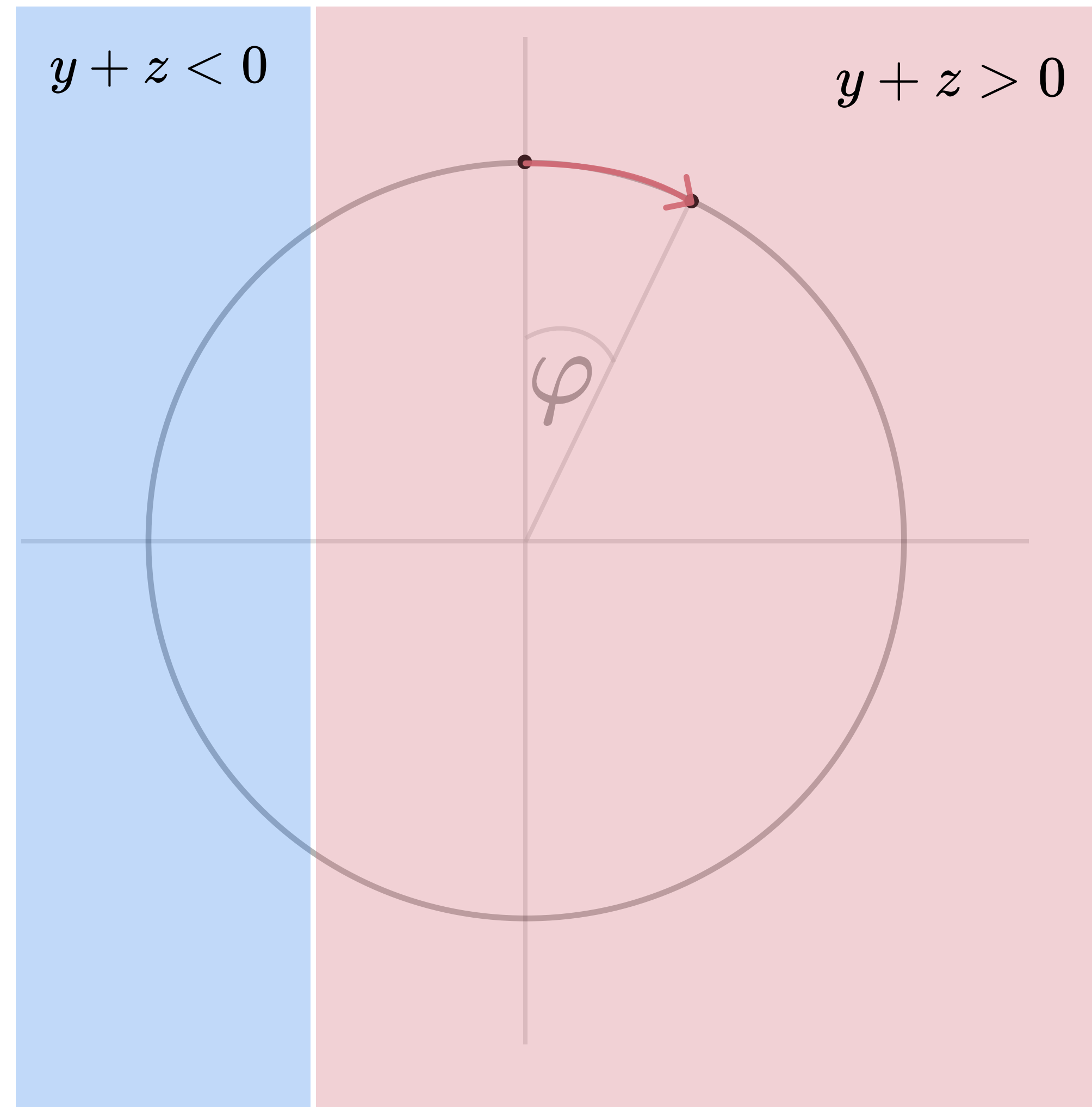



```

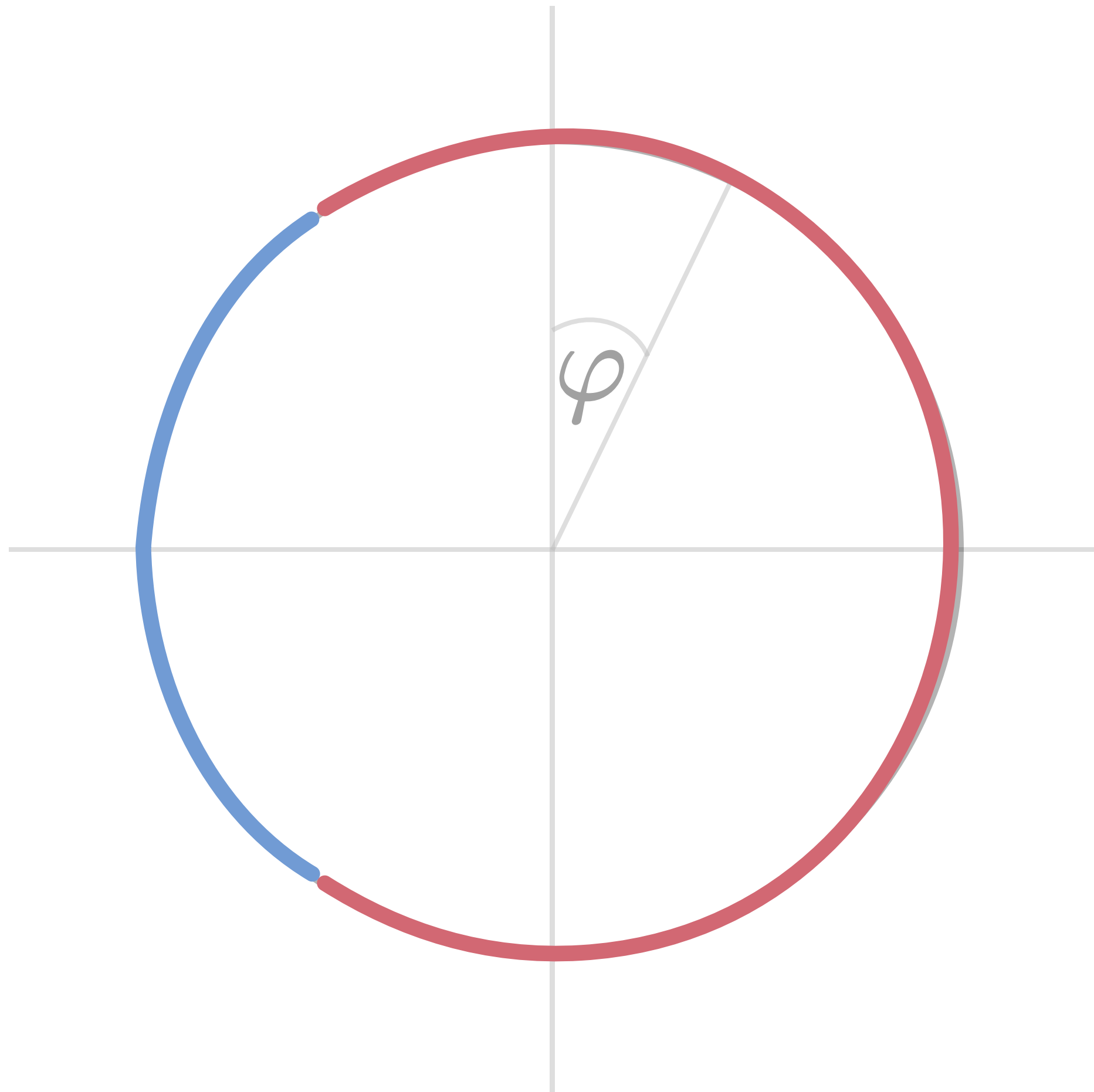
 $x \leftarrow 0; y \leftarrow 6; z \leftarrow 4$ 
while true do
   $x \leftarrow 4x + 3y$ 
   $y \leftarrow 4y - 3x$ 
   $z \leftarrow 5z$ 
  if  $y + z > 0$  then
    Region A
  else
    Region B
  end if
end while

```

Rotation in the first two coordinates by $\varphi = \cos^{-1} 4/5$

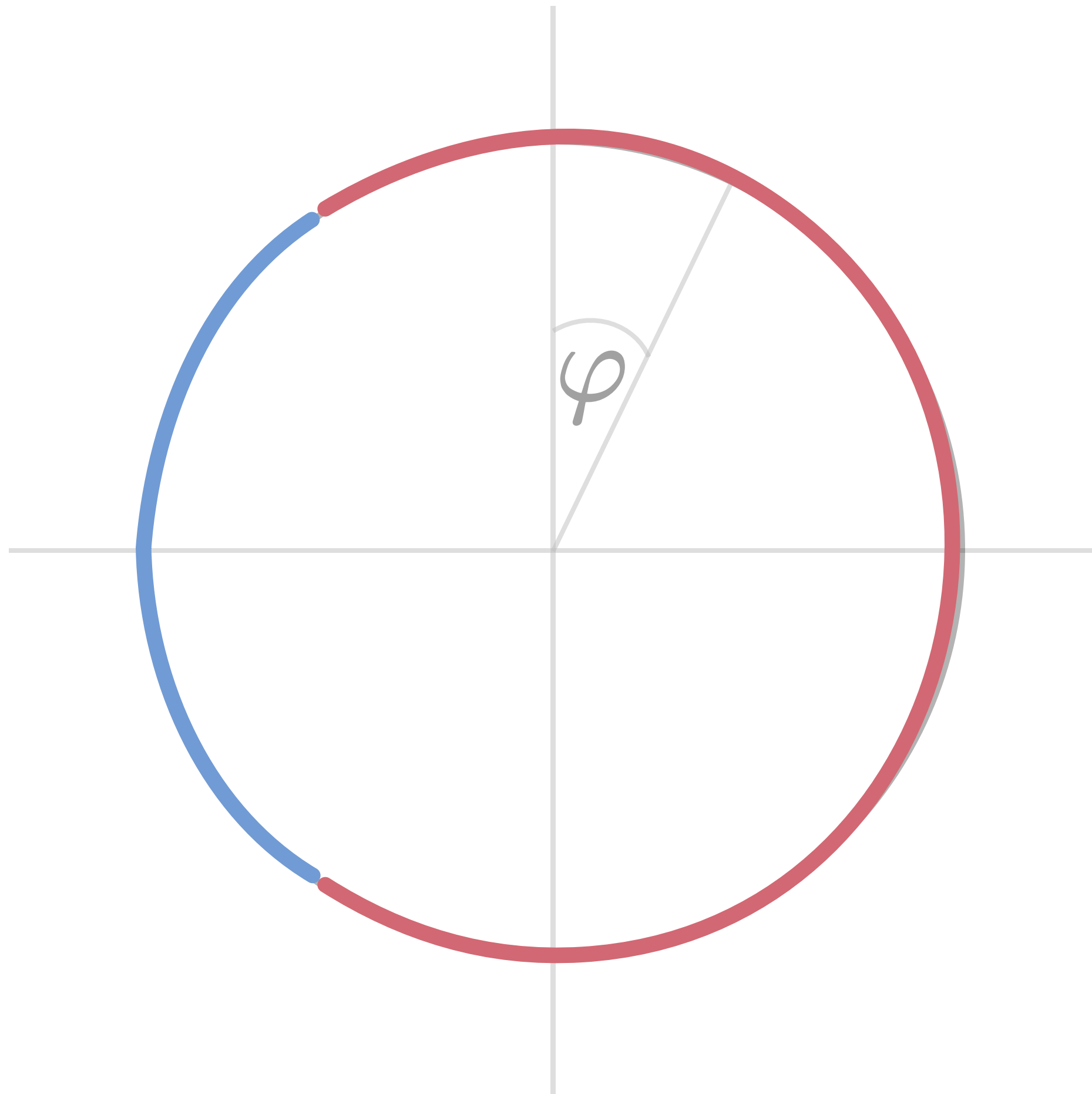


$$\varphi = \cos^{-1} 4/5$$



How frequently are we on the red arc?

$$\varphi = \cos^{-1} 4/5$$



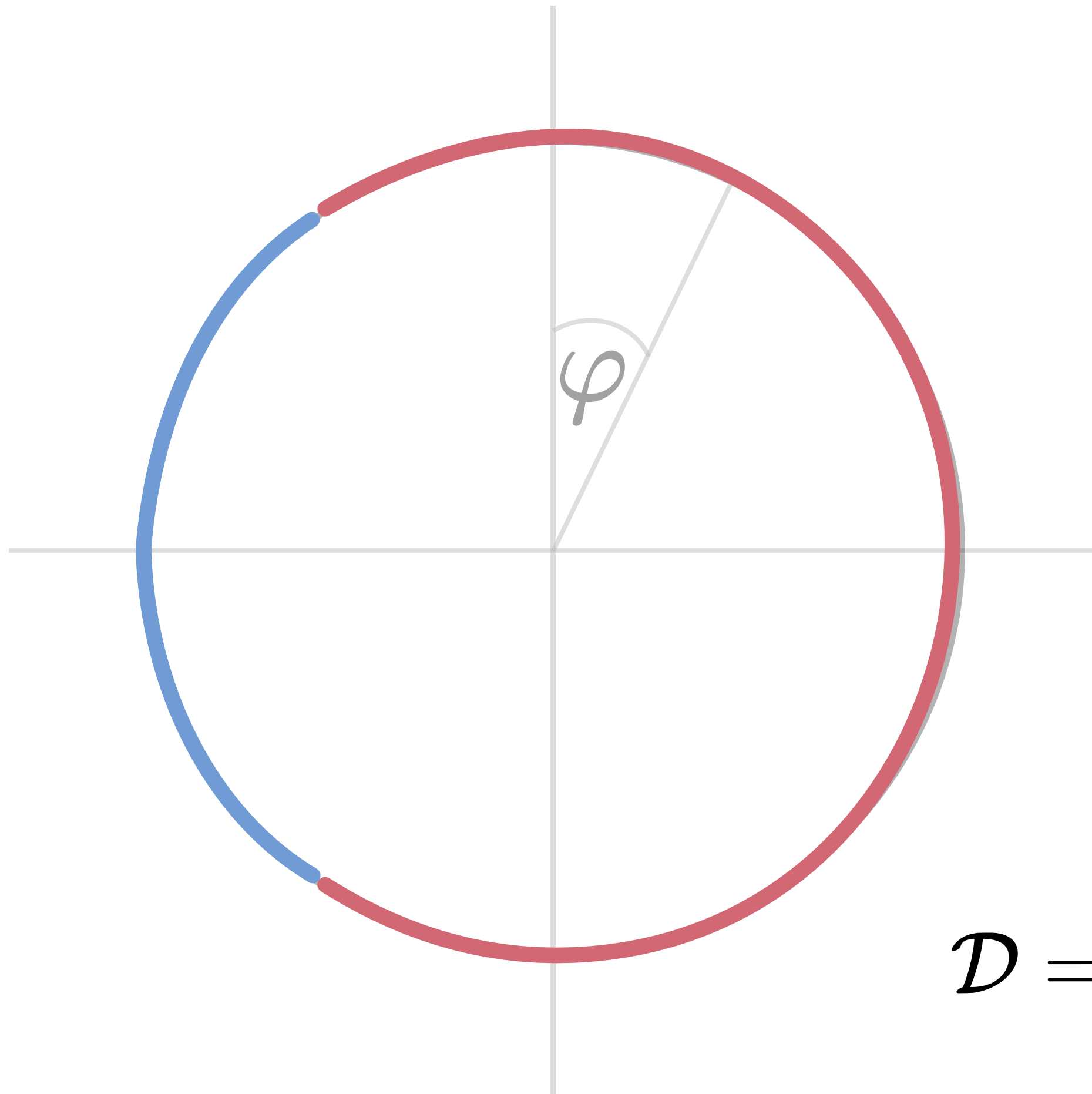
How frequently are we on the red arc?

Theorem (Weyl, 1910). Let ρ be an irrational real number. Then the sequence:

$$\rho, 2\rho, 3\rho, \dots$$

is uniformly distributed mod 1.

$$\varphi = \cos^{-1} 4/5$$



How frequently are we on the red arc?

Theorem (Weyl, 1910). Let ρ be an irrational real number. Then the sequence:

$$\rho, 2\rho, 3\rho, \dots$$

is uniformly distributed mod 1.

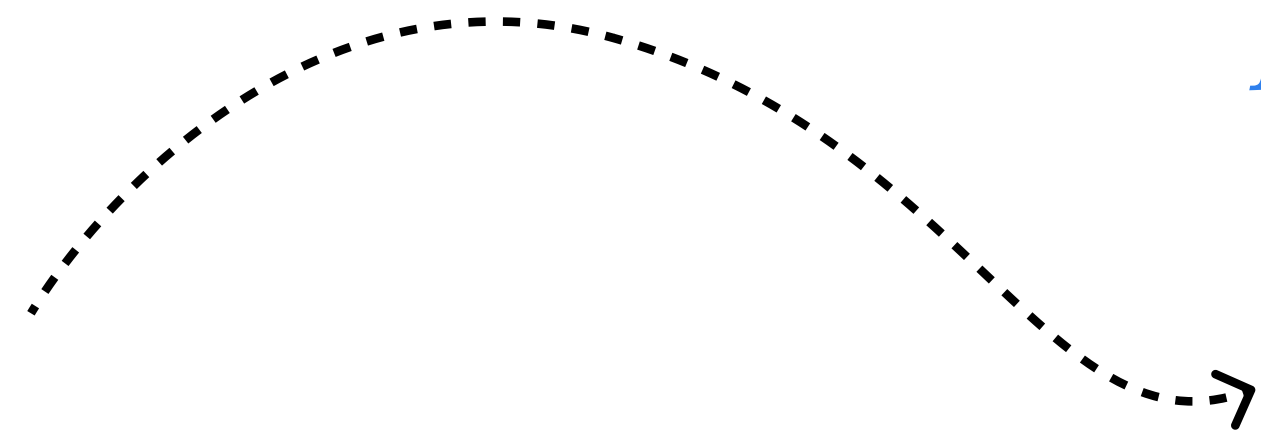
$$\mathcal{D} = \frac{\text{length of } \text{—} \text{—}}{2\pi} = \frac{\cos^{-1}(-2/3)}{\pi} = 0.732278\dots$$

Scheme

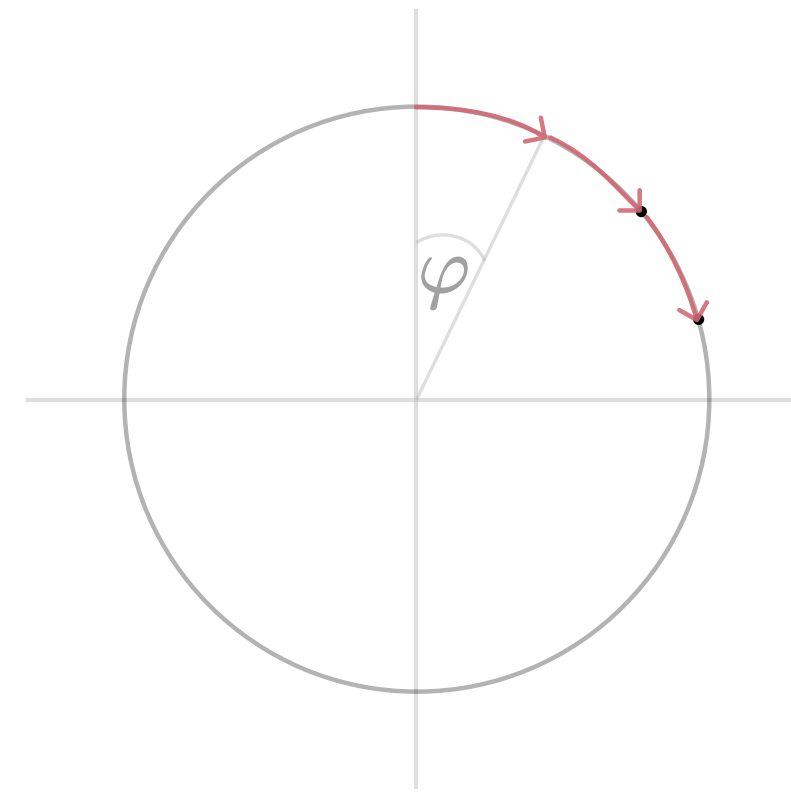
$$\begin{pmatrix} 4 & -3 & 0 \\ 3 & 4 & 0 \\ 0 & 0 & 1 \end{pmatrix}^n$$

Update Matrix

Conserve sign information

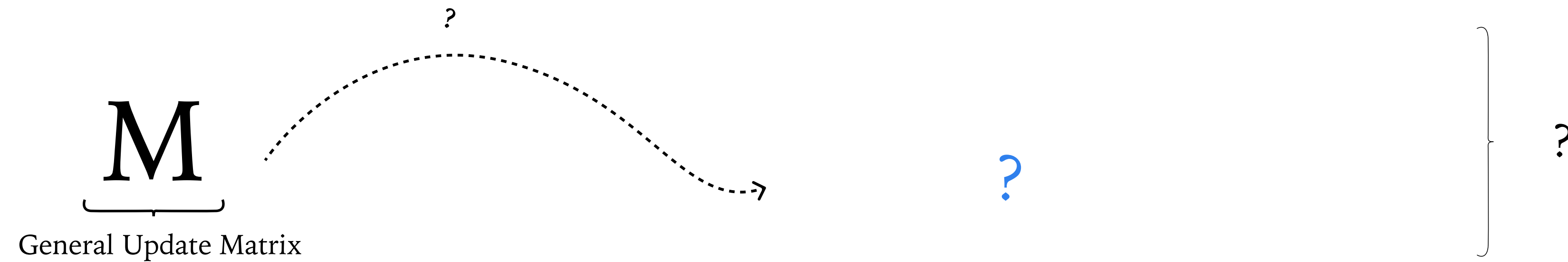
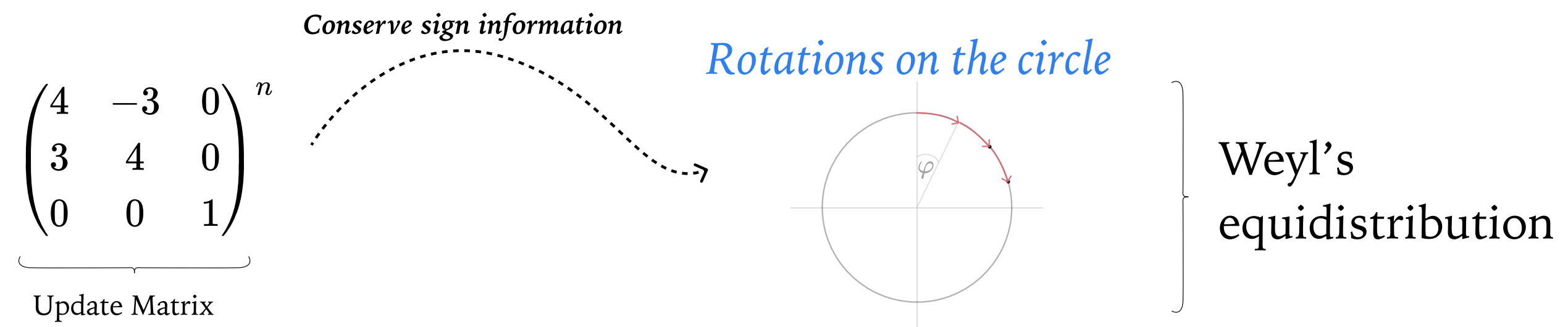


Rotations on the circle

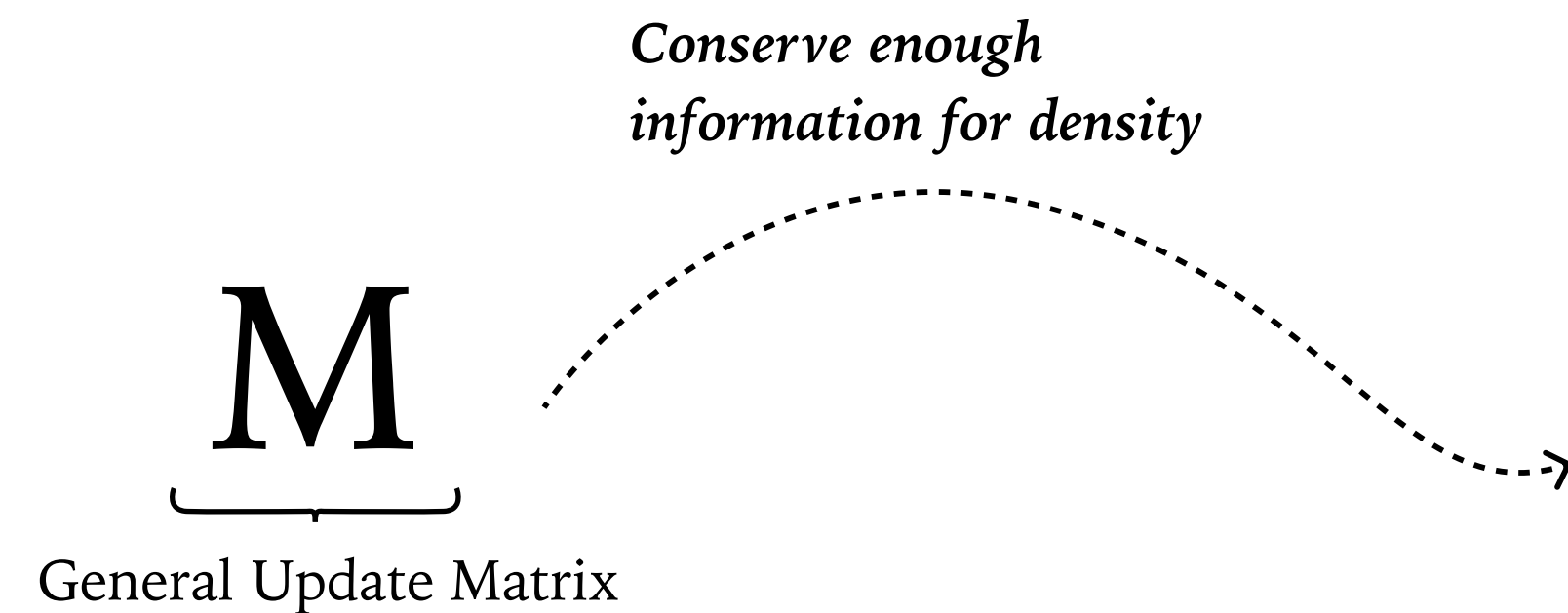
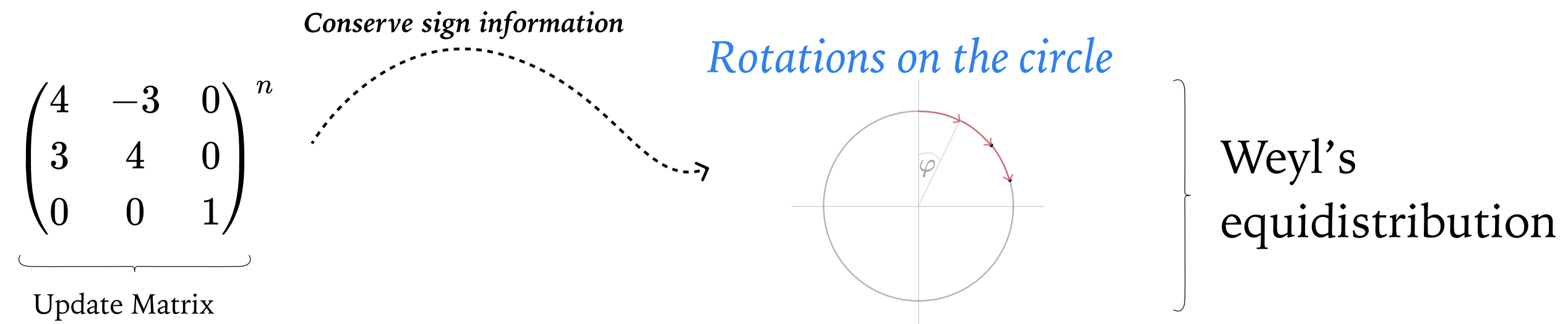


Weyl's
equidistribution

Scheme



Scheme



Rotations on a subgroup of \mathbb{T}^n

A stronger version of Weyl's theorem
+
Koiran's approximation of volumes of definable sets

IV. The Proof

Linear Recurrence Sequences

Theorem (Cayley-Hamilton). $f(\underbrace{M}_{\text{matrix}}) = \mathbf{0}$ characteristic polynomial of M

Linear Recurrence Sequences

Theorem (Cayley-Hamilton). $f(\underbrace{M}_{\text{matrix}}) = \mathbf{0}$ characteristic polynomial of M

$$2M^3 - 4M^2 + M + 5I = \mathbf{0}$$

Linear Recurrence Sequences

Theorem (Cayley-Hamilton). $f(\underbrace{M}_{\text{matrix}}) = \mathbf{0}$ characteristic polynomial of M

$$2M^{n+3} - 4M^{n+2} + M^{n+1} + 5M^n = \mathbf{0}$$

Linear Recurrence Sequences

Theorem (Cayley-Hamilton). $f(\underbrace{M}_{\text{matrix}}) = \mathbf{0}$ characteristic polynomial of M

$$2M^{n+3} - 4M^{n+2} + M^{n+1} + 5M^n = \mathbf{0}$$

$$u_n := (M^n)_{i,j}$$

$$2u_{n+3} - 4u_{n+2} + u_{n+1} + 5u_n = 0$$

Entries of M^n are LRS.

We are interested in the density of $\{n : u_n > 0\}$

Linear Recurrence Sequences

$$\begin{bmatrix} \boxed{\begin{matrix} \lambda_1 & 1 \\ & \lambda_1 & \\ & & \lambda_1 \end{matrix}} & & & \\ & \boxed{\begin{matrix} \lambda_2 & 1 \\ & \lambda_2 \end{matrix}} & & \\ & & \lambda_3 & \\ & & & \dots \\ & & & & \boxed{\begin{matrix} \lambda_n & 1 \\ & \lambda_n \end{matrix}} \end{bmatrix}$$

By Jordan decomposition

$$u_n = \sum_{i=1}^d P_i(n) \lambda_i^n$$

$\in \bar{\mathbb{Q}}$

$\in \bar{\mathbb{Q}}[x]$

Linear Recurrence Sequences

$$\begin{bmatrix} \boxed{\begin{matrix} \lambda_1 & 1 \\ & \lambda_1 & \\ & & \lambda_1 \end{matrix}} & & & \\ & \boxed{\begin{matrix} \lambda_2 & 1 \\ & \lambda_2 \end{matrix}} & & \\ & & \lambda_3 & \\ & & & \ddots \\ & & & & \boxed{\begin{matrix} \lambda_n & 1 \\ & \lambda_n \end{matrix}} \end{bmatrix}$$

By Jordan decomposition

$$u_n = \sum_{i=1}^d P_i(n) \lambda_i^n$$

$\in \overline{\mathbb{Q}}$

$\in \overline{\mathbb{Q}}[x]$

$$\mathcal{D} := \text{density of } \underbrace{\{n : u_n > 0\}}$$

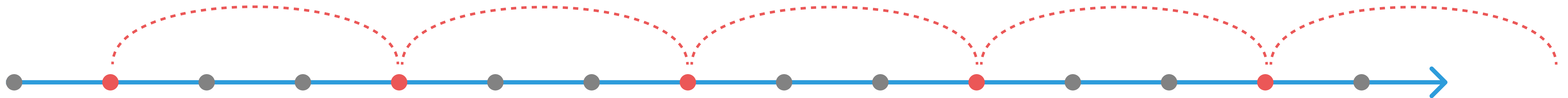
positivity set of $(u_n)_{n \in \mathbb{N}}$

Preprocessing

Split the problem: $(u_{nT+l})_{n \in \mathbb{N}}$, $l \in \{0, 1, \dots, T-1\}$

compute some large period T

where the smaller problems (i.e. subsequences)
have some good properties



Preprocessing

To understand the *good property* of the subsequences, note:

$$u_n = \sum_{i=1}^d P_i(n) \lambda_i^n$$

sometimes there may be multiplicative relations among λ :
integers z_1, \dots, z_d such that:

$$\lambda_1^{z_1} \lambda_2^{z_2} \cdots \lambda_d^{z_d} = 1$$

so e.g. $\lambda_d^{z_d}$ can be written as a product of integer powers of other roots

In the subsequences all the dependencies are gathered

Preprocessing

Good property: **the subsequences have the same sign as**

$$\underbrace{\sum_{i \in I} c_i \alpha_i^n}_{\text{independent}} + \underbrace{\sum_{i \in D} c_i \prod_{j \in I} \alpha_j^{q_{i,j}}}_{\text{dependent}} + \underbrace{c + R(n)}_{\text{remainder tends to zero}}, \quad n \in \mathbb{N}$$

- $\alpha_i \in \bar{\mathbb{Q}}, \quad |\alpha_i| = 1$

Furthermore, the group

- $c_i \in \bar{\mathbb{Q}}$

$$\{(z_1, \dots, z_k) \in \mathbb{Z}^k : \underbrace{\alpha_1^{z_1} \alpha_2^{z_2} \cdots \alpha_k^{z_k}}_{\text{independent}} = 1\} = \{\mathbf{0}\}$$

- $q_{i,j} \in \mathbb{Q}$

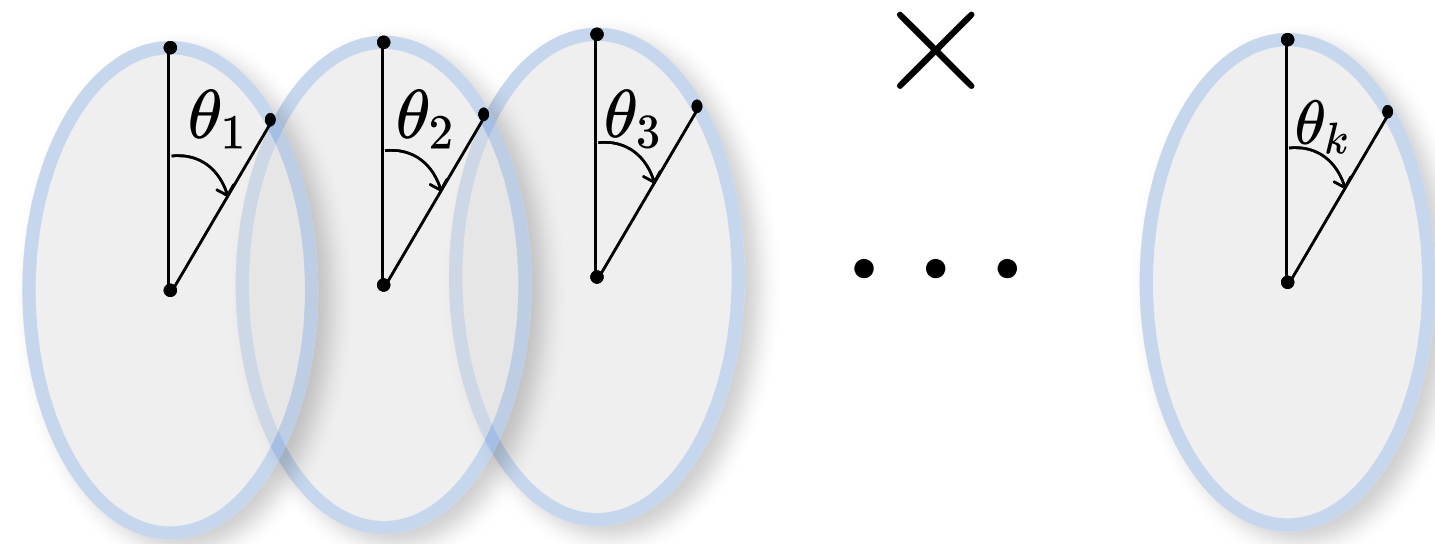
is trivial.

Rotations

$$\alpha_i = e^{2\pi i \theta_i}$$

$$(\alpha_1^n, \alpha_2^n, \dots, \alpha_k^n), \quad n \in \mathbb{N}$$

$\{\theta_1, \dots, \theta_k, 1\}$
are linearly
independent over \mathbb{Q}

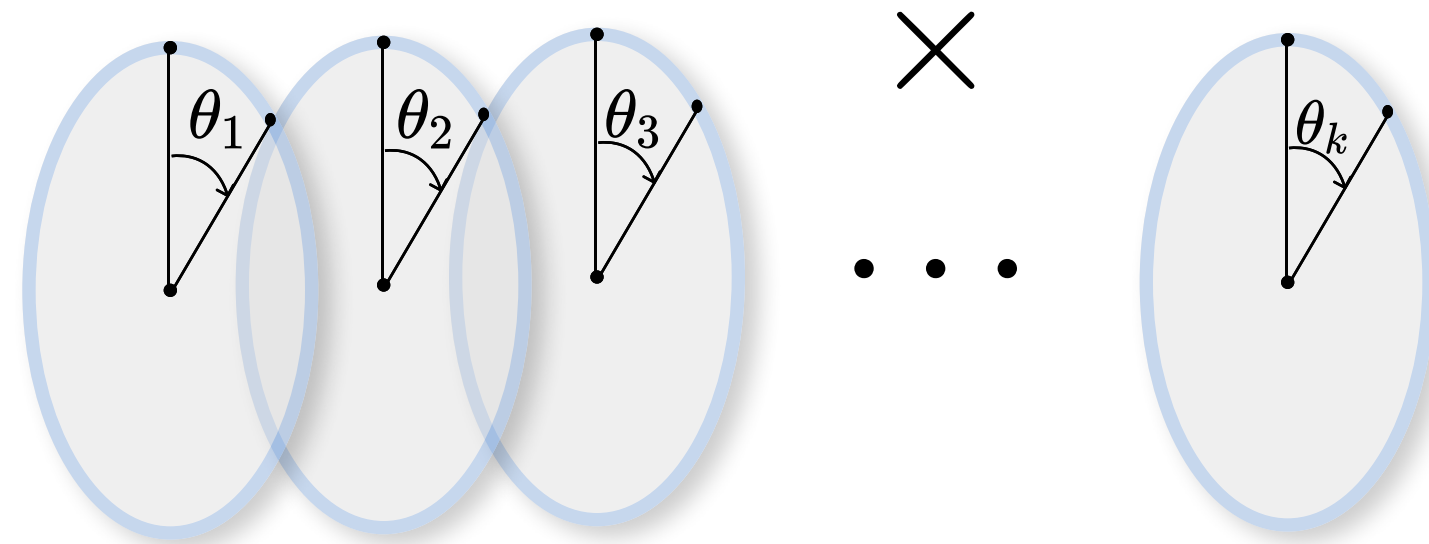


Rotations

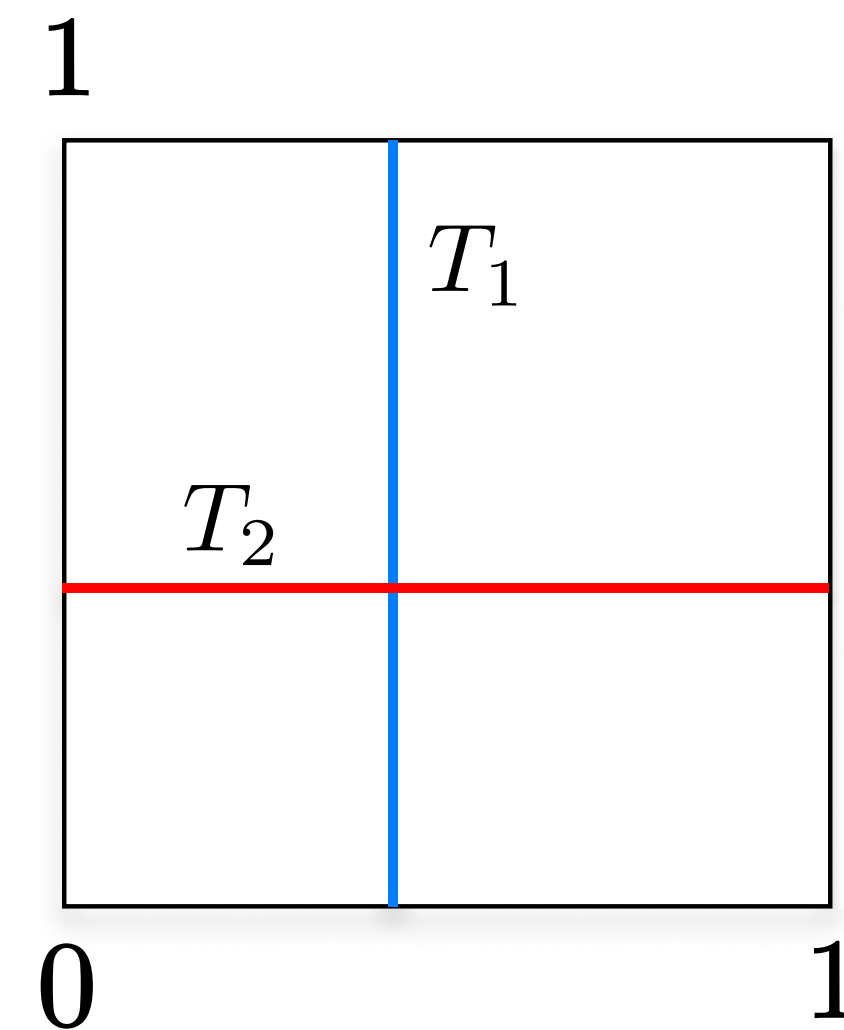
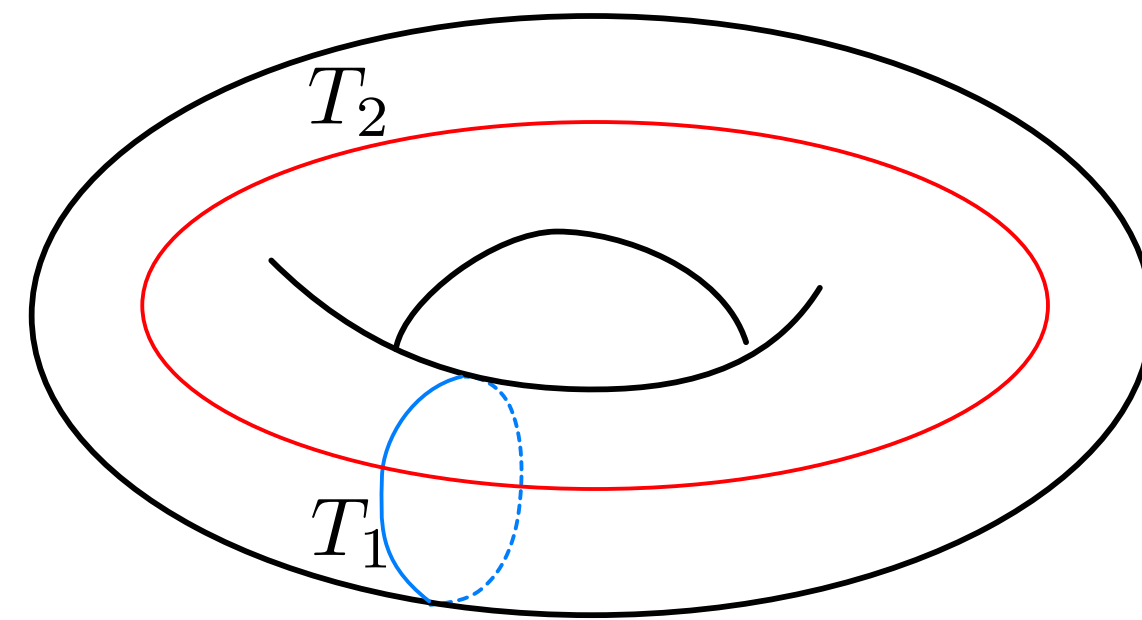
$$\alpha_i = e^{2\pi i \theta_i}$$

$$(\alpha_1^n, \alpha_2^n, \dots, \alpha_k^n), \quad n \in \mathbb{N}$$

$\{\theta_1, \dots, \theta_k, 1\}$
are linearly
independent over \mathbb{Q}



E.g. for $k=2$



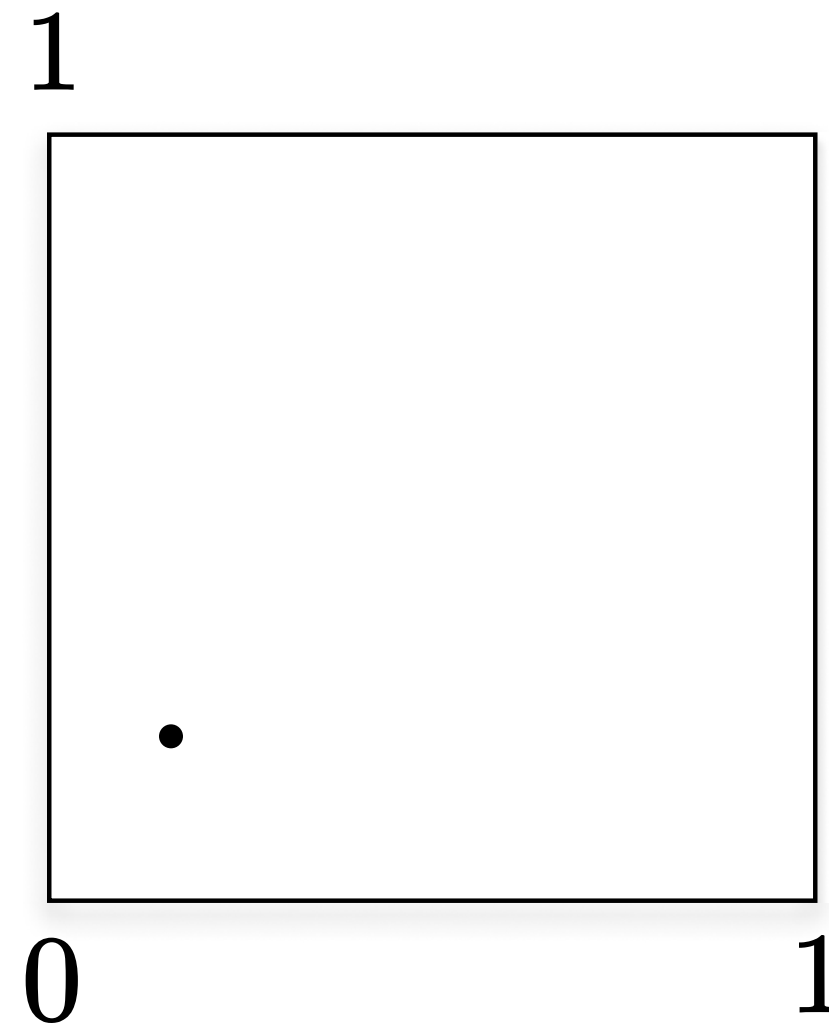
Rotations

$$\alpha_i = e^{2\pi i \theta_i}$$

$$(\alpha_1^n, \alpha_2^n, \dots, \alpha_k^n), \quad n \in \mathbb{N}$$

$\{\theta_1, \dots, \theta_k, 1\}$
are linearly
independent over \mathbb{Q}

$$x_i + \theta_i \pmod{1}$$



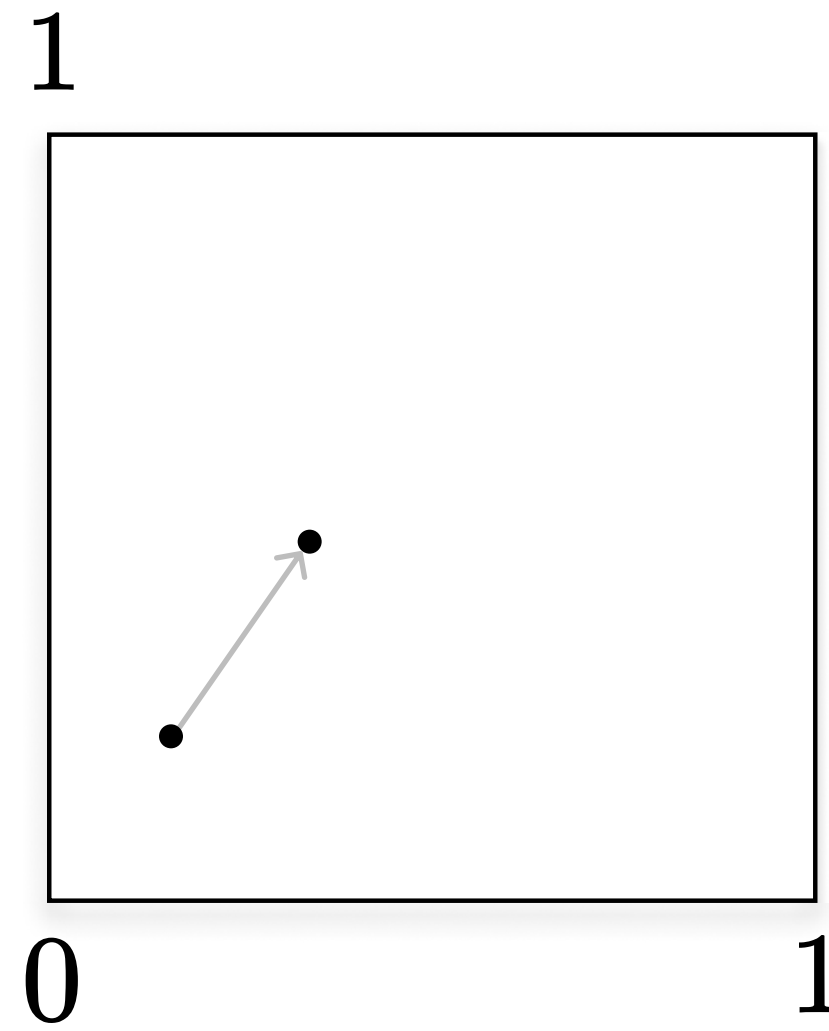
Rotations

$$\alpha_i = e^{2\pi i \theta_i}$$

$$(\alpha_1^n, \alpha_2^n, \dots, \alpha_k^n), \quad n \in \mathbb{N}$$

$\{\theta_1, \dots, \theta_k, 1\}$
are linearly
independent over \mathbb{Q}

$$x_i + \theta_i \pmod{1}$$



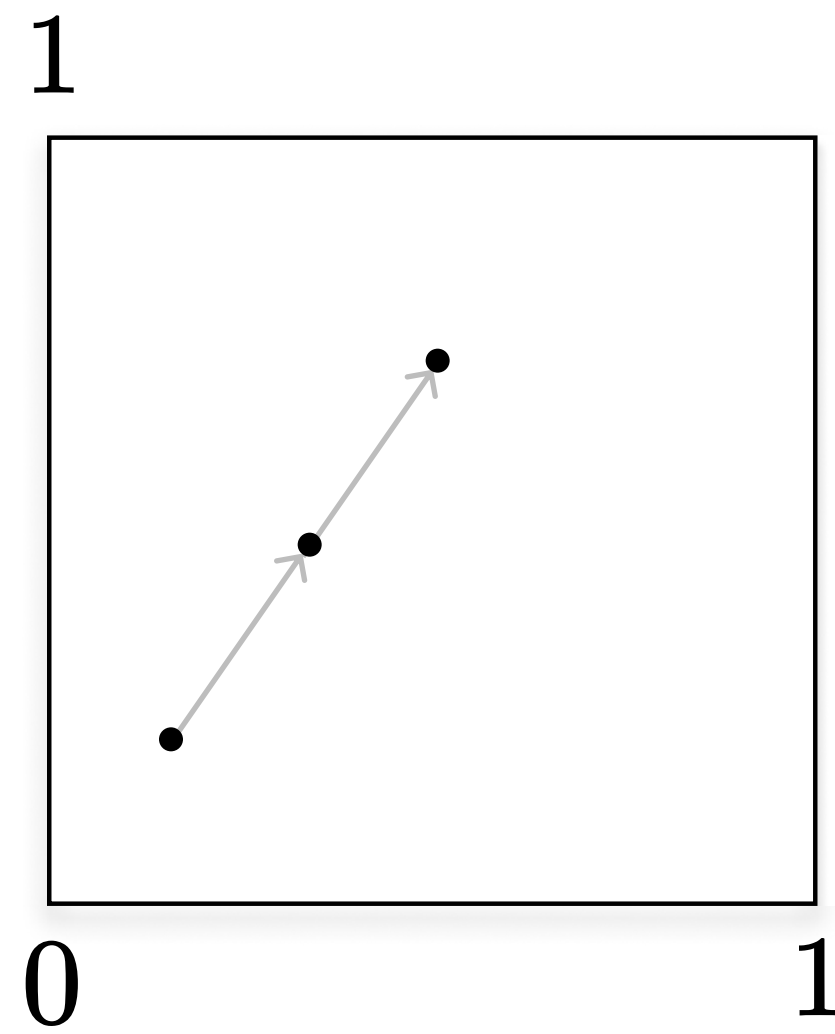
Rotations

$$\alpha_i = e^{2\pi i \theta_i}$$

$$(\alpha_1^n, \alpha_2^n, \dots, \alpha_k^n), \quad n \in \mathbb{N}$$

$\{\theta_1, \dots, \theta_k, 1\}$
are linearly
independent over \mathbb{Q}

$$x_i + \theta_i \pmod{1}$$



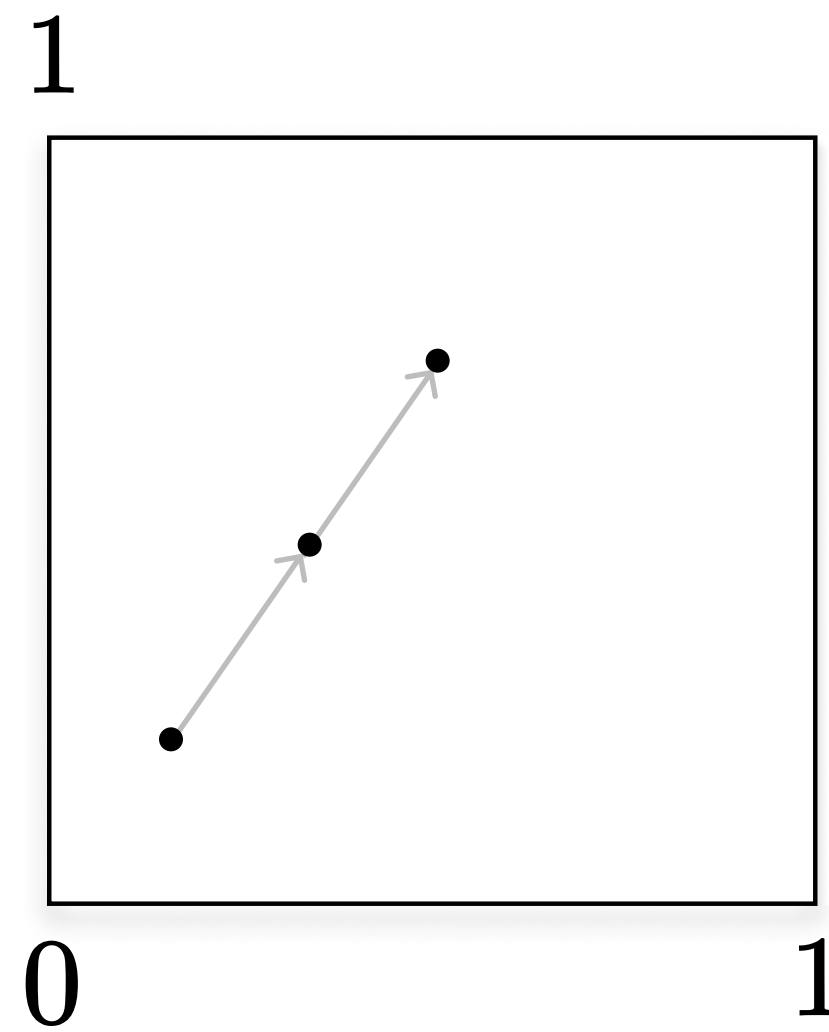
Rotations

$$\alpha_i = e^{2\pi i \theta_i}$$

$$(\alpha_1^n, \alpha_2^n, \dots, \alpha_k^n), \quad n \in \mathbb{N}$$

$\{\theta_1, \dots, \theta_k, 1\}$
are linearly
independent over \mathbb{Q}

$$x_i + \theta_i \pmod{1}$$



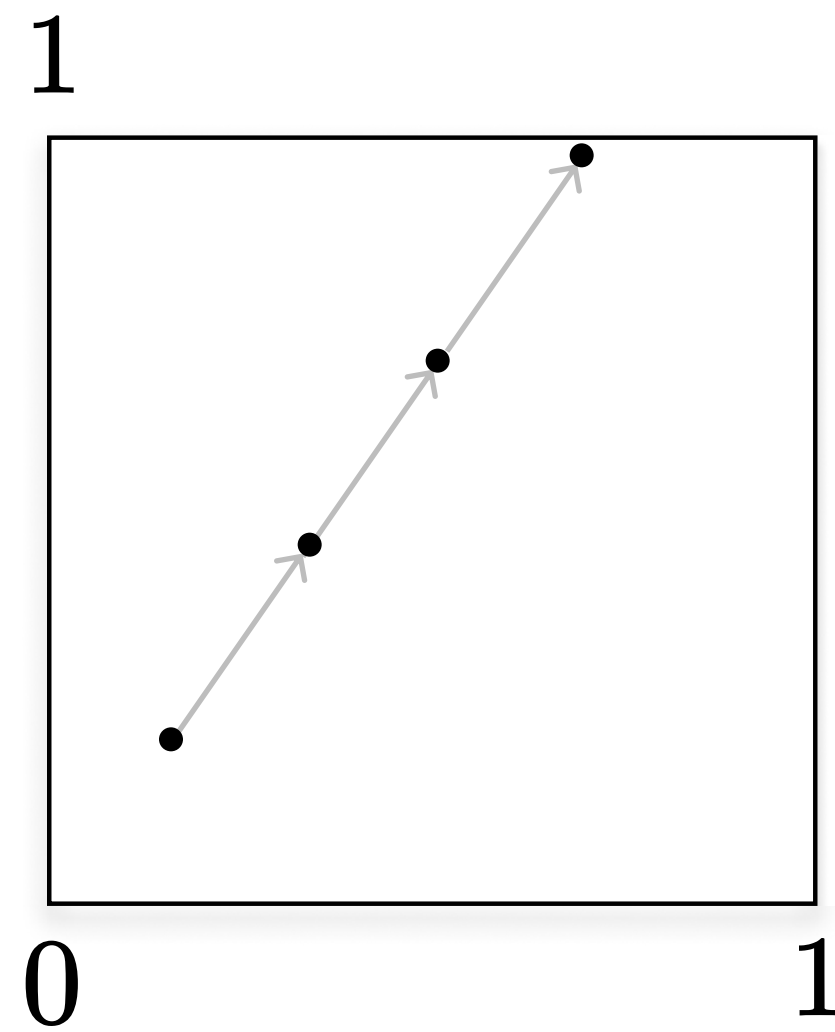
Rotations

$$\alpha_i = e^{2\pi i \theta_i}$$

$$(\alpha_1^n, \alpha_2^n, \dots, \alpha_k^n), \quad n \in \mathbb{N}$$

$\{\theta_1, \dots, \theta_k, 1\}$
are linearly
independent over \mathbb{Q}

$$x_i + \theta_i \pmod{1}$$



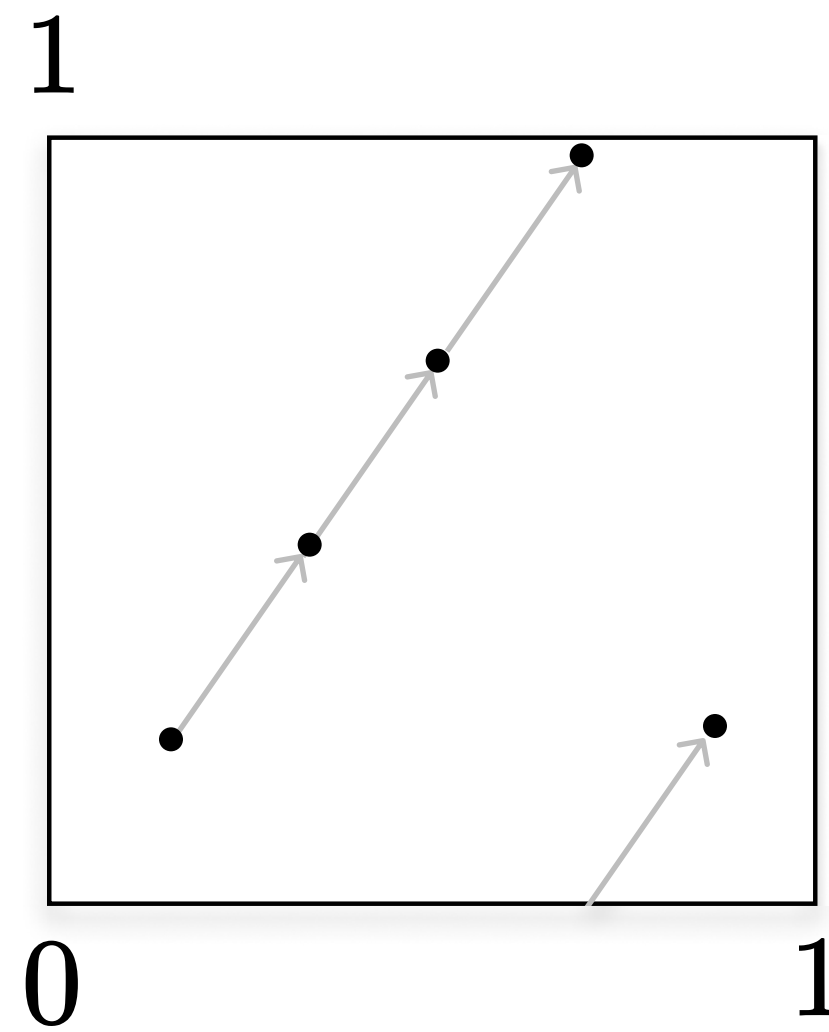
Rotations

$$\alpha_i = e^{2\pi i \theta_i}$$

$$(\alpha_1^n, \alpha_2^n, \dots, \alpha_k^n), \quad n \in \mathbb{N}$$

$\{\theta_1, \dots, \theta_k, 1\}$
are linearly
independent over \mathbb{Q}

$$x_i + \theta_i \pmod{1}$$



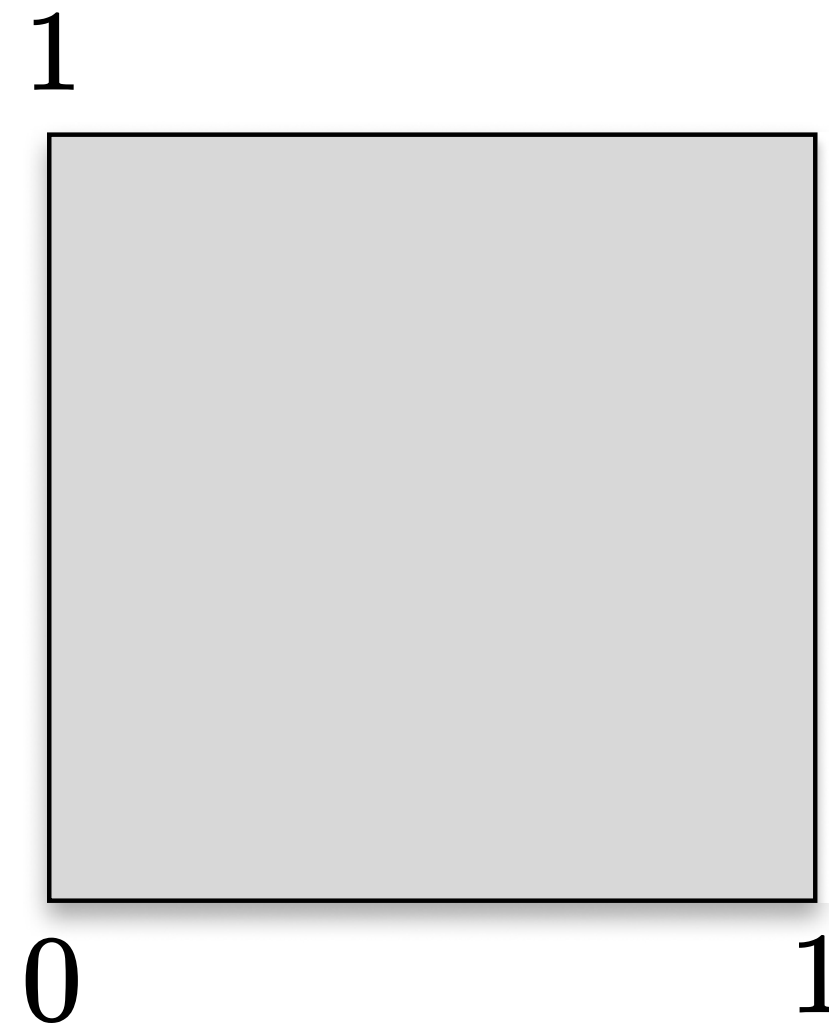
Rotations

$$\alpha_i = e^{2\pi i \theta_i}$$

$$(\alpha_1^n, \alpha_2^n, \dots, \alpha_k^n), \quad n \in \mathbb{N}$$

$\{\theta_1, \dots, \theta_k, 1\}$
are linearly
independent over \mathbb{Q}

$$x_i + \theta_i \pmod{1}$$



Theorem (Kronecker).
 $(n\theta_1 \pmod{1}, \dots, n\theta_k \pmod{1}), n \in \mathbb{N}$
is dense in the hypercube

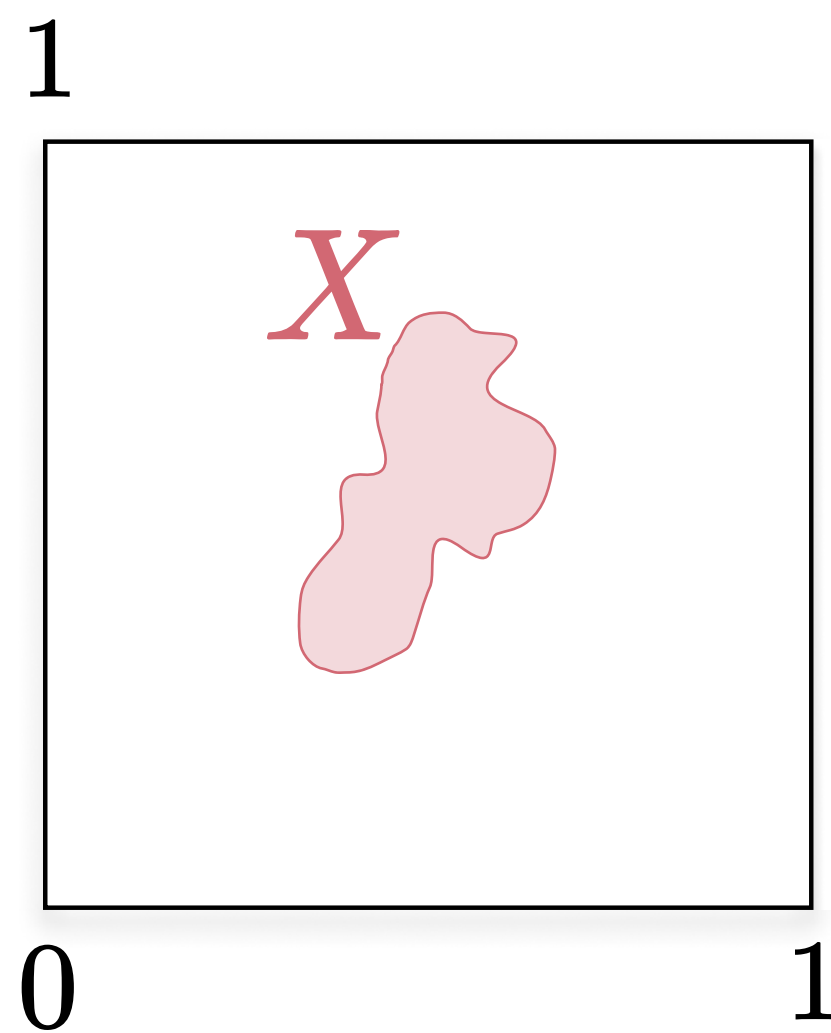
Rotations

$$\alpha_i = e^{2\pi i \theta_i}$$

$$(\alpha_1^n, \alpha_2^n, \dots, \alpha_k^n), \quad n \in \mathbb{N}$$

$\{\theta_1, \dots, \theta_k, 1\}$
are linearly
independent over \mathbb{Q}

$$x_i + \theta_i \pmod{1}$$



Theorem (Weyl, 1912).

$$(n\theta_1 \pmod{1}, \dots, n\theta_k \pmod{1}), \quad n \in \mathbb{N}$$

is *equidistributed* in the
hypercube

The amount of time spent in X is proportional to $\text{vol}(X)$

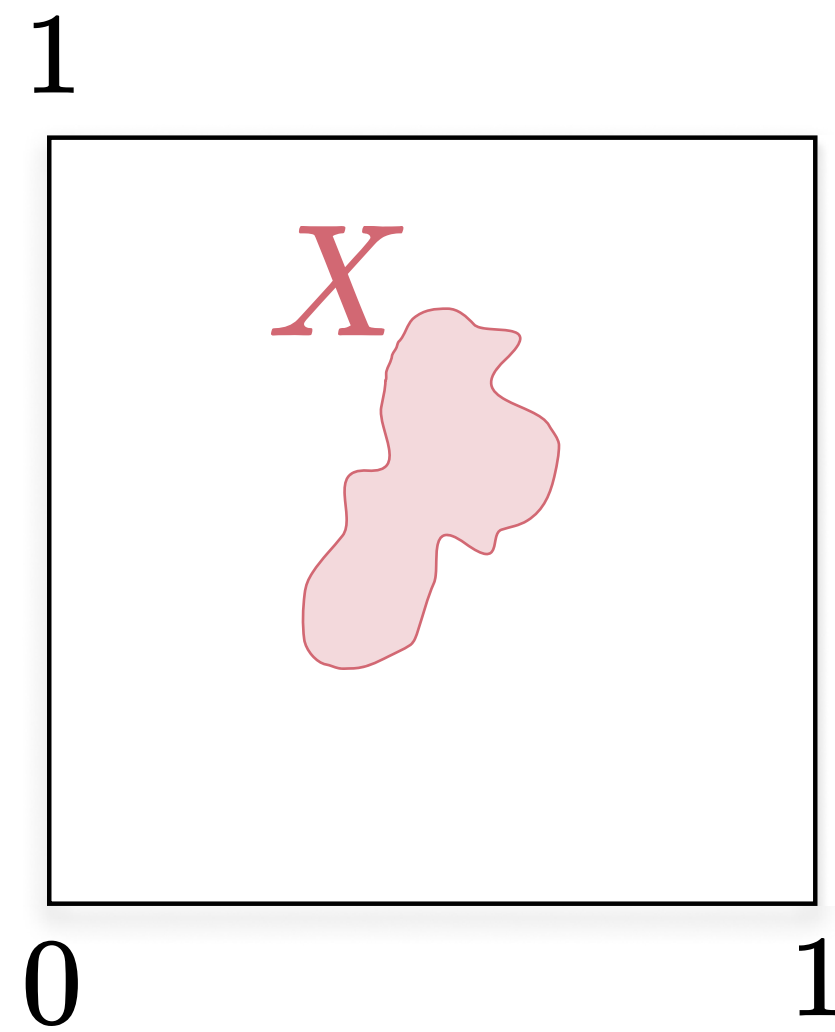
Rotations

$$\alpha_i = e^{2\pi i \theta_i}$$

$$(\alpha_1^n, \alpha_2^n, \dots, \alpha_k^n), \quad n \in \mathbb{N}$$

$\{\theta_1, \dots, \theta_k, 1\}$
are linearly
independent over \mathbb{Q}

$$x_i + \theta_i \pmod{1}$$



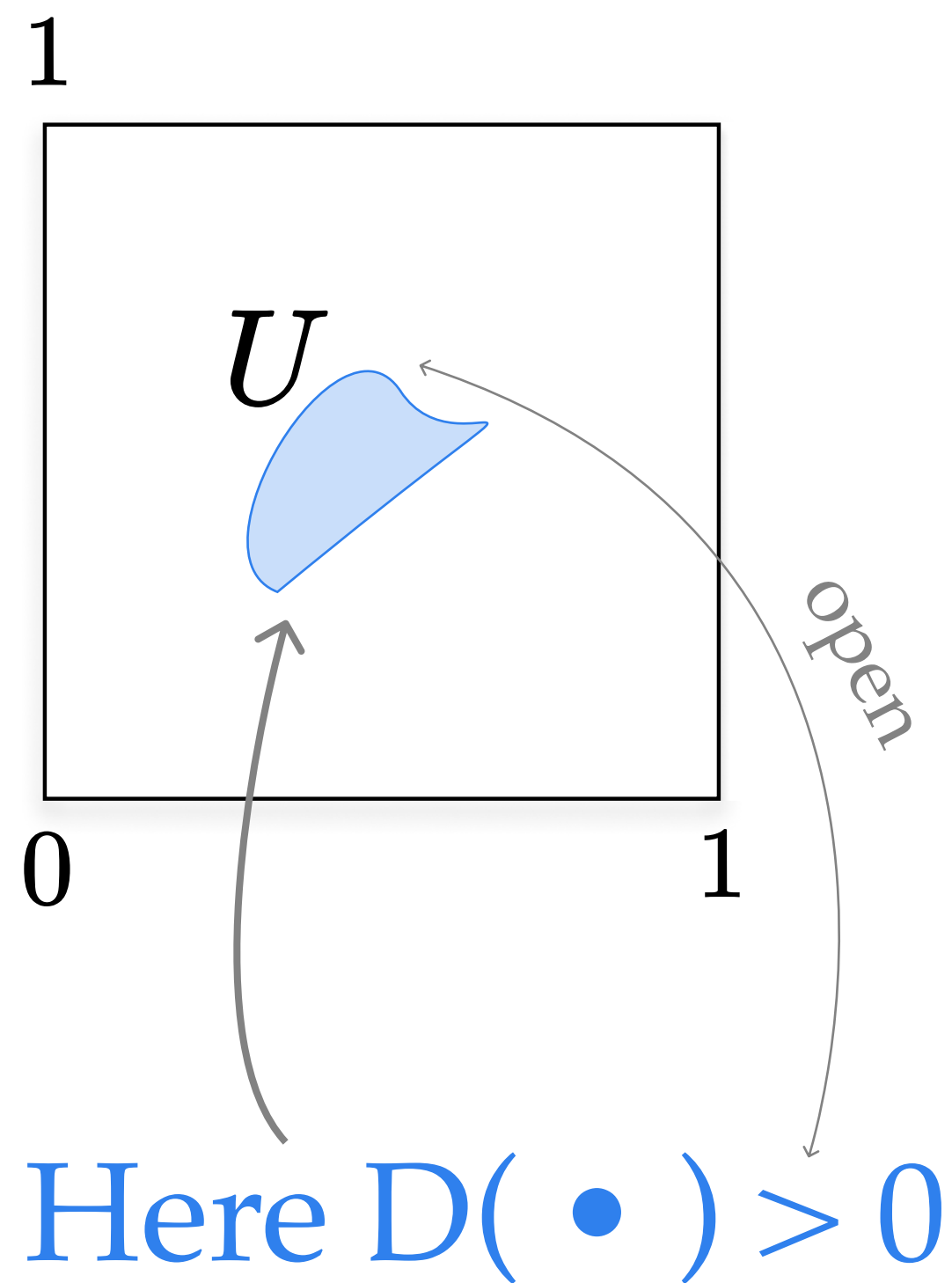
Theorem (Weyl, 1912).
 $(n\theta_1 \pmod{1}, \dots, n\theta_k \pmod{1}), n \in \mathbb{N}$
is *equidistributed* in the
hypercube

The amount of time spent in X is proportional to $\text{vol}(X)$

$$\text{Density of } \{n : (n\theta_1 \pmod{1}, \dots, n\theta_k \pmod{1}) \in X\} = \text{vol}(X)$$

$$\underbrace{\sum_{i \in I} c_i \alpha_i^n + \sum_{i \in D} c_i \prod_{j \in I} \alpha_j^{q_{i,j}} + c}_{D(n)} + R(n), \quad n \in \mathbb{N}$$

$$\underbrace{\sum_{i \in I} c_i \alpha_i^n + \sum_{i \in D} c_i \prod_{j \in I} \alpha_j^{q_{i,j}} + c}_{D(n)} + R(n), \quad n \in \mathbb{N}$$

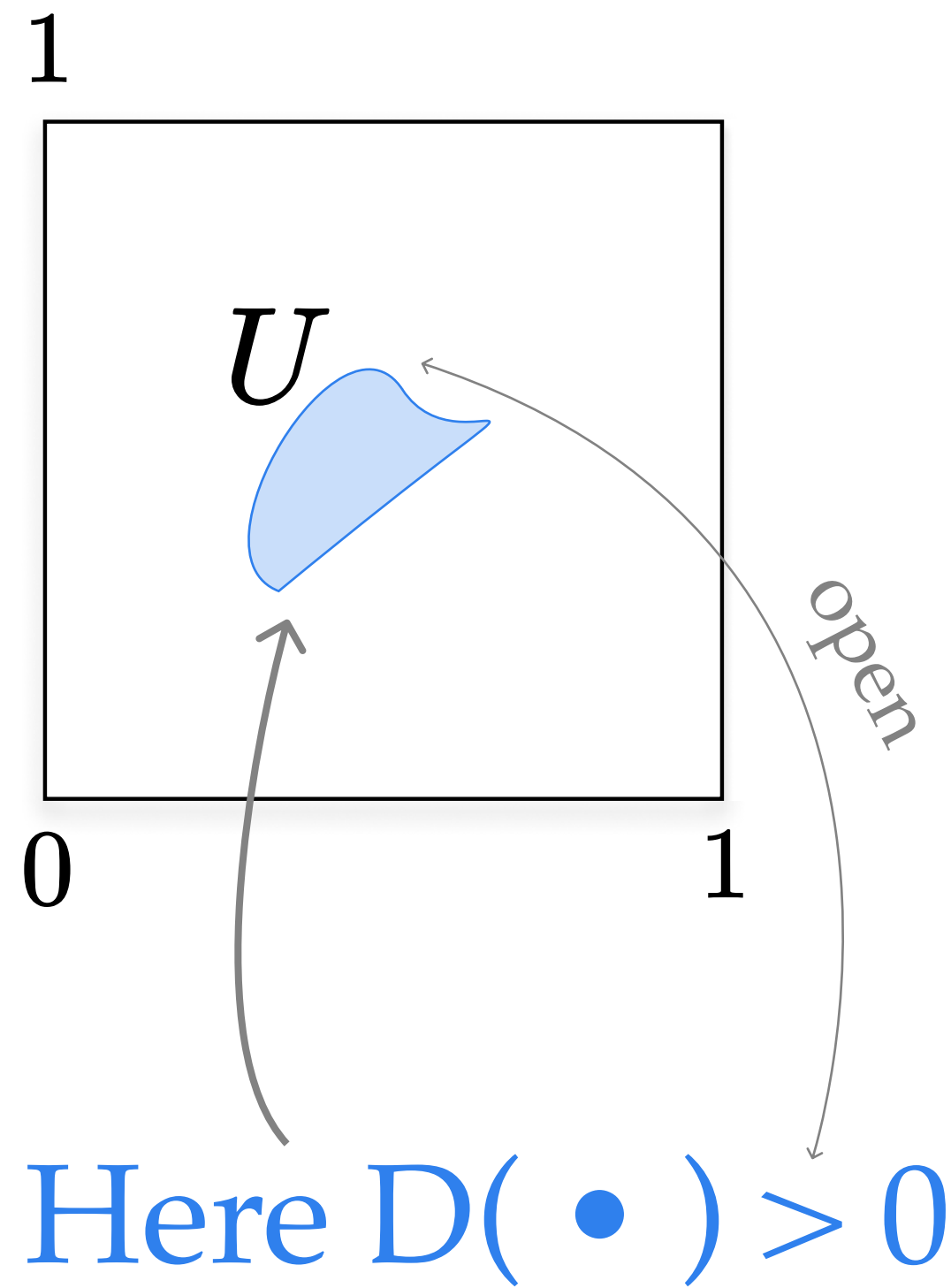


Density of $\{n : D(n) > 0\} = \text{vol}(U)$

$$\text{vol}(U) > 0 \iff \underbrace{U \neq \emptyset}$$

An equivalent statement can be decided with Tarski's algo

$$\underbrace{\sum_{i \in I} c_i \alpha_i^n + \sum_{i \in D} c_i \prod_{j \in I} \alpha_j^{q_{i,j}} + c}_{D(n)} + R(n), \quad n \in \mathbb{N}$$



$$\text{Density of } \{n : D(n) > 0\} = \text{vol}(U)$$

$$\text{vol}(U) > 0 \iff \underbrace{U \neq \emptyset}$$

An equivalent statement can be decided with FO of reals

so we can decide if the density of the positivity set of $D(n)$ is nonzero

$$\underbrace{\sum_{i \in I} c_i \alpha_i^n + \sum_{i \in D} c_i \prod_{j \in I} \alpha_j^{q_{i,j}} + c}_{D(n)} + R(n), \quad n \in \mathbb{N}$$

We need: Density of $\{n : D(n) + R(n) > 0\}$

$$\underbrace{\sum_{i \in I} c_i \alpha_i^n + \sum_{i \in D} c_i \prod_{j \in I} \alpha_j^{q_{i,j}} + c}_{D(n)} + R(n), \quad n \in \mathbb{N}$$

We need: Density of $\{n : D(n) + R(n) > 0\}$

When is $|D(n)| < |R(n)|$?

Difficult problem: Depends on diophantine properties of α
 (Main obstruction to decidability of Skolem, Positivity, etc.)

$$\underbrace{\sum_{i \in I} c_i \alpha_i^n + \sum_{i \in D} c_i \prod_{j \in I} \alpha_j^{q_{i,j}} + c}_{D(n)} + R(n), \quad n \in \mathbb{N}$$

We need: Density of $\{n : D(n) + R(n) > 0\}$

When is $|D(n)| < |R(n)|$?

Difficult problem: Depends on diophantine properties of α
(Main obstruction to decidability of Skolem, Positivity, etc.)

$$\text{Density of } \{n : |D(n)| < |R(n)|\} = 0$$

When is $|D(n)| < |R(n)|$?

Density of $\{n : |D(n)| < |R(n)|\} = 0$

- Density of $\{n : D(n) = 0\} = 0$
due to Skolem-Mahler-Lech
- $\lim_{n \rightarrow \infty} |R(n)| = 0$ polynomially fast

When is $|D(n)| < |R(n)|$?

Density of $\{n : |D(n)| < |R(n)|\} = 0$

- Density of $\{n : D(n) = 0\} = 0$
due to Skolem-Mahler-Lech
- $\lim_{n \rightarrow \infty} |R(n)| = 0$ polynomially fast

Theorem 1. “ $\mathcal{D} = 0?$ ” is decidable.

(so is “ $\mathcal{D} = 1?$ ” by symmetry)

Diagonalisable Matrices

When M is diagonalisable:

$$\lim_{n \rightarrow \infty} |R(n)| = 0 \text{ exponentially fast}$$

$$\underbrace{\sum_{i \in I} c_i \alpha_i^n + \sum_{i \in D} c_i \prod_{j \in I} \alpha_j^{q_{i,j}} + c}_{D(n)} + R(n), \quad n \in \mathbb{N}$$

In this case, the p-adic subspace theorem implies:

not effective

there is some N , such that for all $n > N$

$$|D(n)| > |R(n)|.$$

$D(n) + R(n) > 0$ for infinitely many n
if and only if
 $D(n) > 0$ for infinitely many n

Diagonalisable Matrices

When M is diagonalisable:

$\lim_{n \rightarrow \infty} |R(n)| = 0$ exponentially fast

$$\underbrace{\sum_{i \in I} c_i \alpha_i^n + \sum_{i \in D} c_i \prod_{j \in I} \alpha_j^{q_{i,j}} + c}_{D(n)} + R(n), \quad n \in \mathbb{N}$$

In this case, the p-adic subspace theorem implies:

there is some N , such that for all $n > N$
 $|D(n)| > |R(n)|$.

not effective

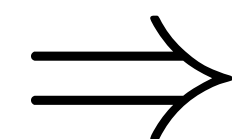
Theorem 1a does not hold for nondiagonalisable matrices.

Joël Ouaknine , James Worrell :

Ultimate Positivity is Decidable for Simple Linear Recurrence Sequences. ICALP (2) 2014: 330-341

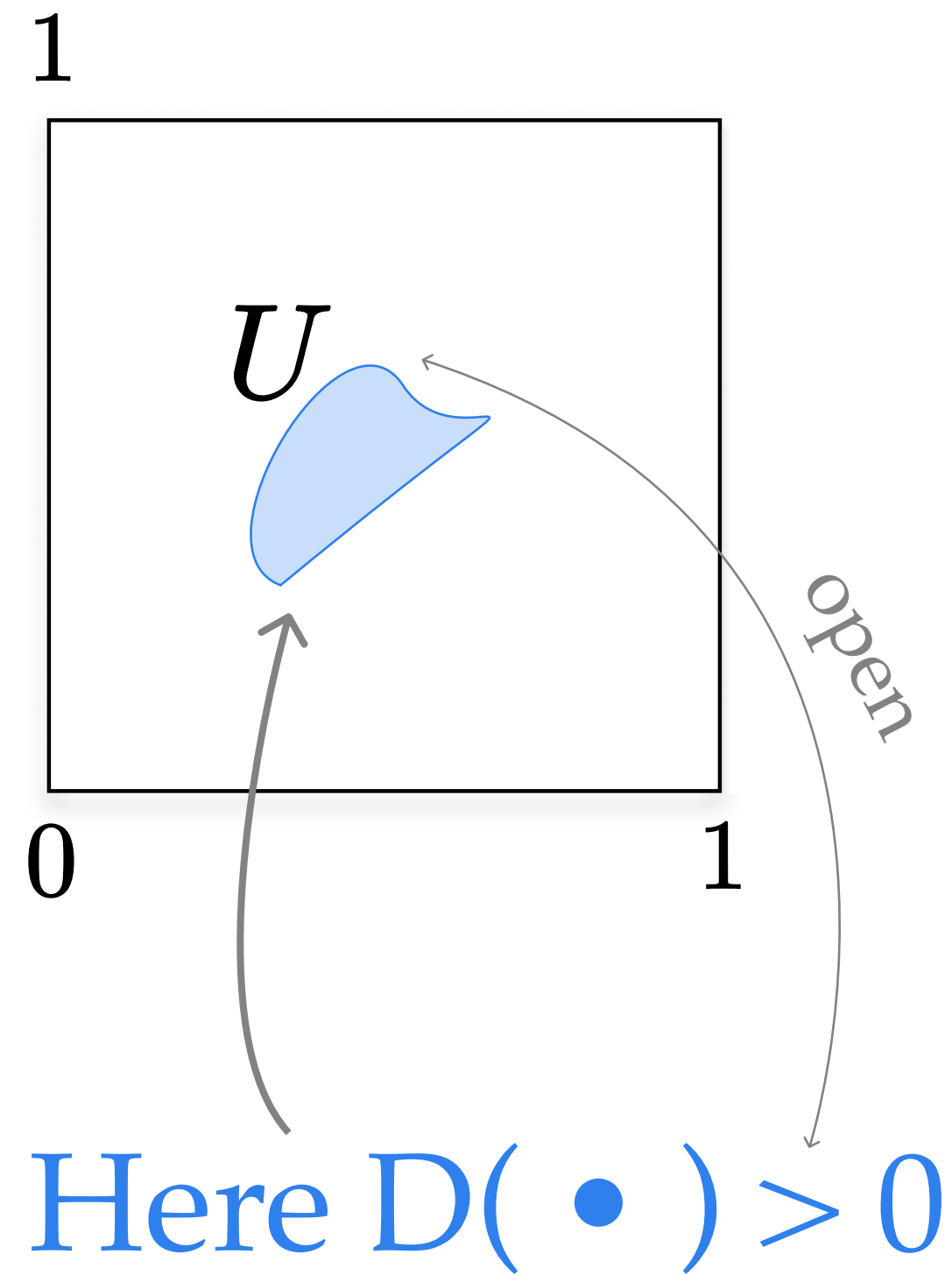
$D(n) + R(n) > 0$ for infinitely many n
 if and only if

$D(n) > 0$ for infinitely many n



Theorem 1a. For diagonalisable update matrices:
 $\mathcal{D} = 0 \Leftrightarrow$ finitely many ■

Approximating the Density

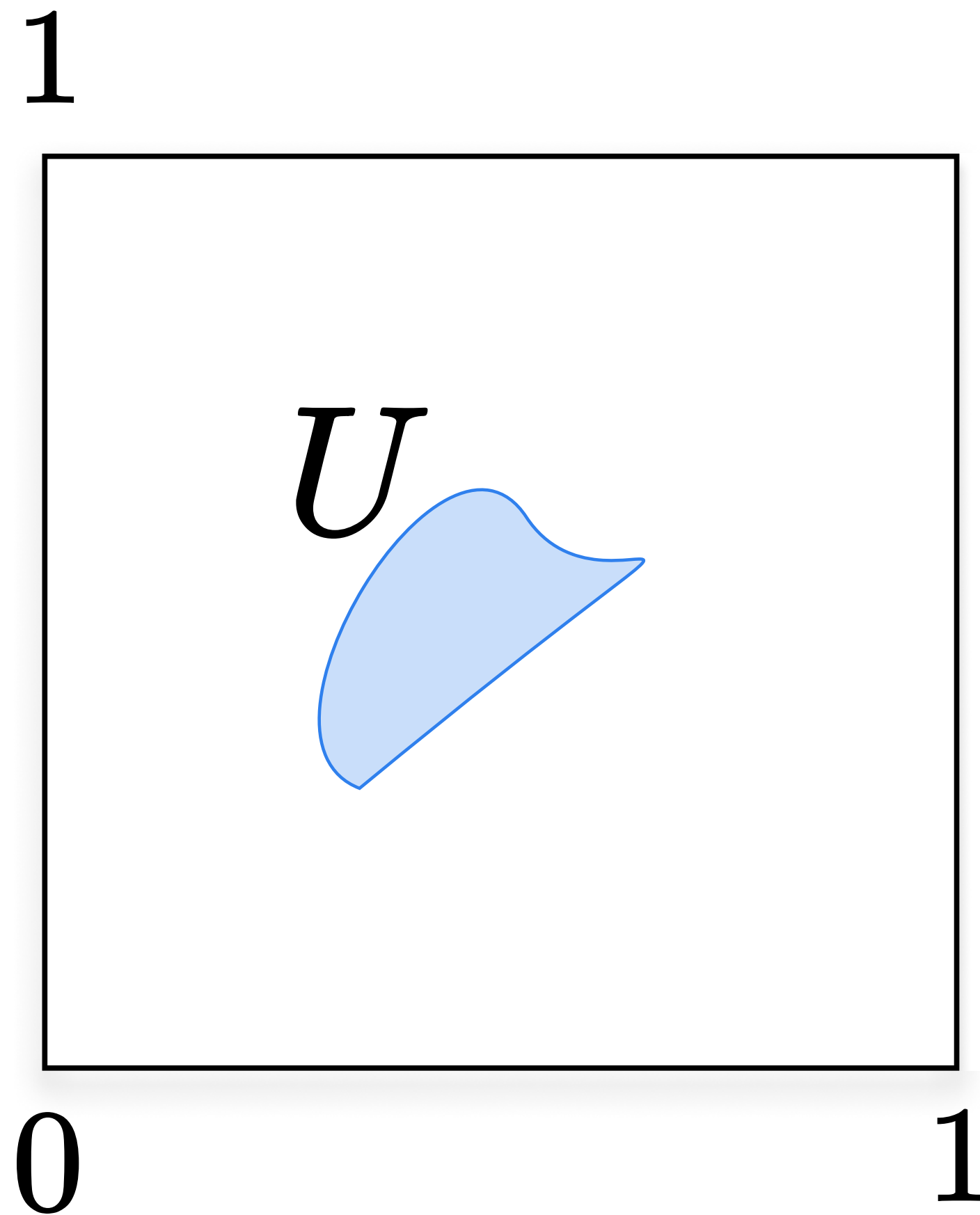


can be even transcendental

$$\text{Density of } \{n : D(n) > 0\} = \text{vol}(U)$$

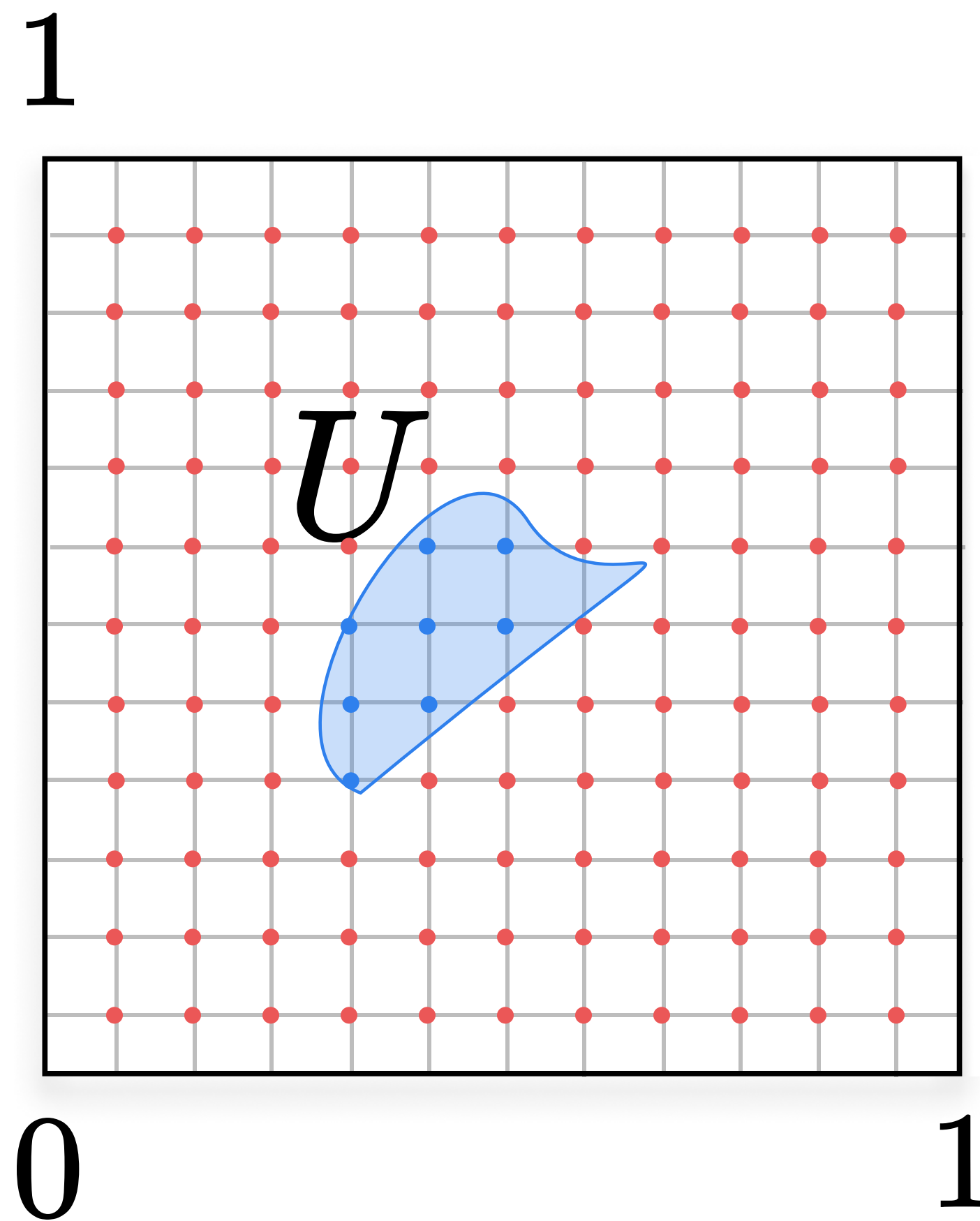
Approximating the Density

How to approximate the volume of U ?



Approximating the Density

How to approximate the volume of U ?

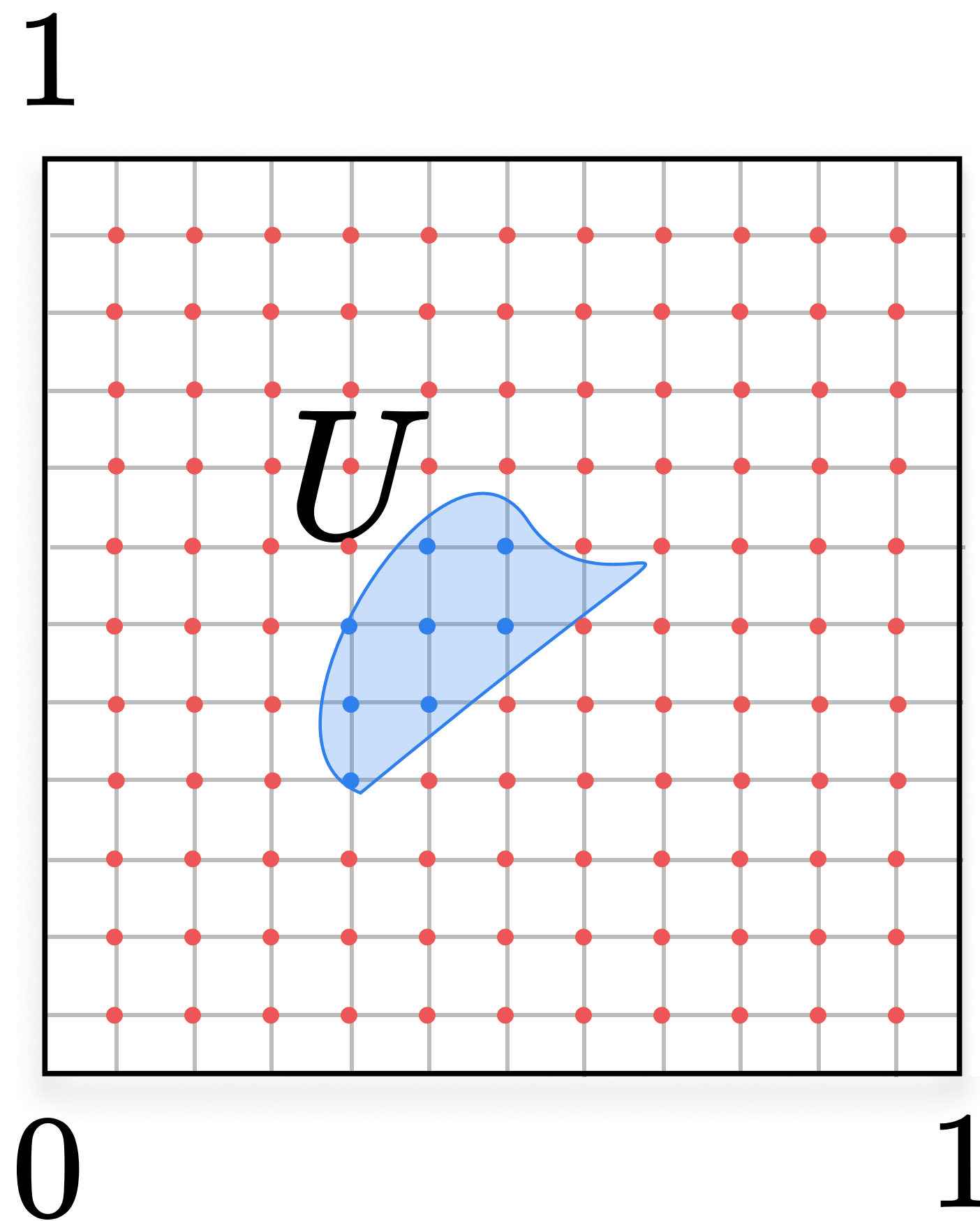


Approximate volume
 \approx
 $\frac{\text{number of } \bullet}{\text{number of } \bullet}$

Approximating the Density

How to approximate the volume of U ?

1. How to make a grid such that $\bullet \in U$ can be decided
2. How fine should the grid be for $|\text{approx} - \text{vol}| < \varepsilon$?



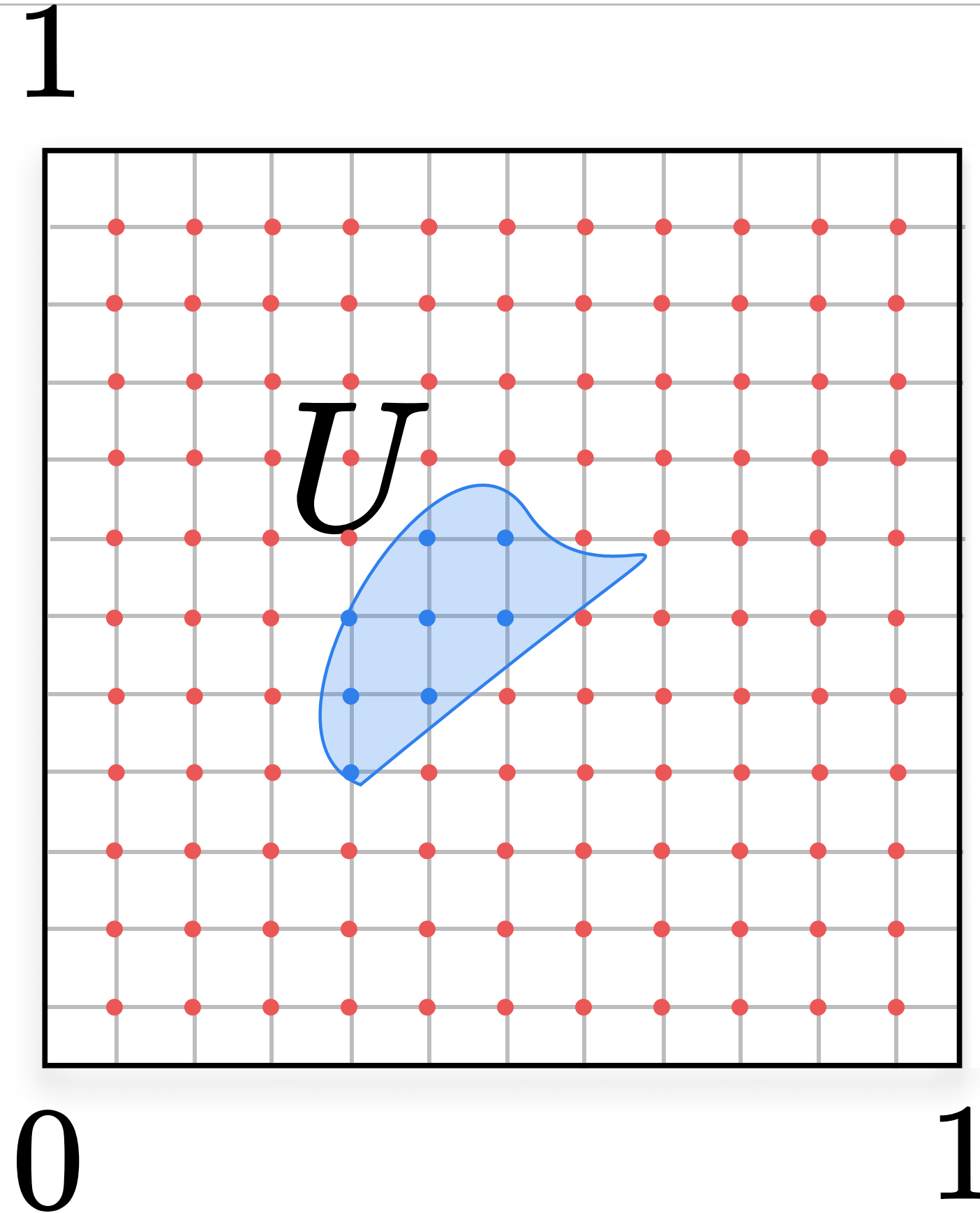
$$\text{Approximate volume} \approx \frac{\text{number of } \bullet}{\text{number of } \bullet}$$

Approximating the Density

1. How to make a grid such that $\bullet \in U$ can be decided
2. How fine should the grid be for $|\text{approx} - \text{vol}| < \varepsilon$?

Pascal Koiran:
Approximating the Volume of Definable Sets. FOCS 1995: 134-141

there is also a Monte-Carlo
type algorithm



Approximate volume
 \parallel
 $\frac{\text{number of } \bullet}{\text{number of } \bullet}$

Theorem 2. \mathcal{D} can be computed to arbitrary precision.

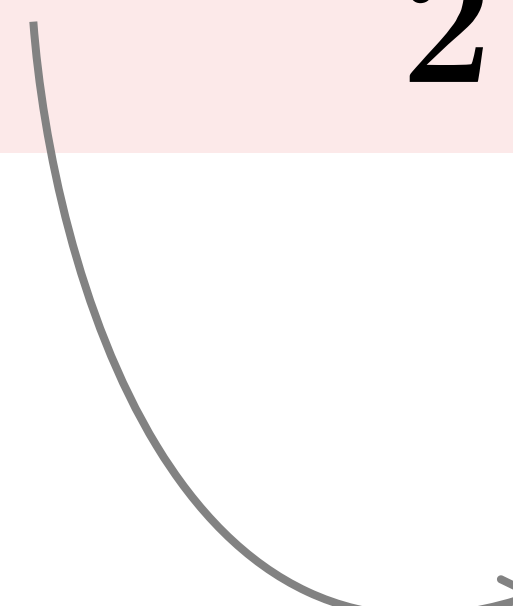
Complexity

- **< PSPACE**
- when number of variables (order of LRS) is fixed, **< PTIME**
- **> co-NP**

V. The Open Problem

Decide whether

$$\mathcal{D} > \frac{1}{2}$$



Can be rational, algebraic, or
transcendental

Decide whether

$$\mathcal{D} > \frac{1}{2}$$

If $\mathcal{D} \notin \mathbb{Q}$ we can use the approximation algorithm

Can be rational, algebraic, or transcendental

If $\mathcal{D} \in \mathbb{Q}$ then we can probably compute it directly

Can we decide whether $\mathcal{D} \in \mathbb{Q}$?

Decide whether $\mathcal{D} \in \mathbb{Q}$

When there are no multiplicative relations:

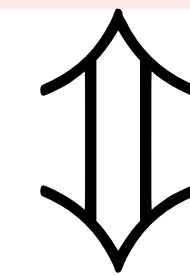
Decide if $\int_{\mathcal{L}} \prod_{i=1}^{\eta} \frac{1}{\sqrt{1-x_i^2}} d\vec{x} \in \mathbb{Q} \pi^{\eta}$

polytope 

Decide whether $\mathcal{D} \in \mathbb{Q}$

When there are at most three dominant eigenvalues,
the problem reduces to:

Given $\alpha \in \overline{\mathbb{Q}}$, decide whether $\cos^{-1}(\alpha) \in \mathbb{Q}\pi$



$$T_n(\cos \theta) = \cos(n\theta)$$

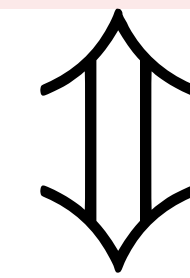
for some n , α is a root of

$$T_n(x) - 1 \quad \text{or} \quad T_n(x) + 1$$

Decide whether $\mathcal{D} \in \mathbb{Q}$

When there are at most three dominant eigenvalues,
the problem reduces to:

Given $\alpha \in \bar{\mathbb{Q}}$, decide whether $\cos^{-1}(\alpha) \in \mathbb{Q}\pi$



$$T_n(\cos \theta) = \cos(n\theta)$$

for some n , α is a root of

$$T_n(x) - 1 \quad \text{or} \quad T_n(x) + 1$$

Theorem 3. “ $\mathcal{D} \in \mathbb{Q}$?” is decidable,
when there are at most three dominant eigenvalues.

Thank You

I. The Problem

```

x ← 0; y ← 6; z ← 4
while true do
  x ← 4x + 3y
  y ← 4y - 3x
  z ← 5z
  if y + z > 0 then
    Region A
  else
    Region B
  end if
end while
    
```



Decision questions:

1. Is **Region A** reached?
(Is there at least one ■?)
 - Known as the **positivity problem**;
 - at least as hard as Skolem's problem
2. Is **Region A** reached infinitely often?
(Are there infinitely many ■?)
 - Known as the **ultimate positivity problem**;
 - also open & difficult

In this paper:

3. How much more frequent are ■ compared to ■?

Set of ■

1. Is it empty?
2. Is it infinite?
3. How big is it inside \mathbb{N} ?

II. The Theorems

Theorems

Theorem 1. “ $\mathcal{D} = 0$?” is decidable.

(so is “ $\mathcal{D} = 1$?” by symmetry)

Theorem 1a. For diagonalisable update matrices:

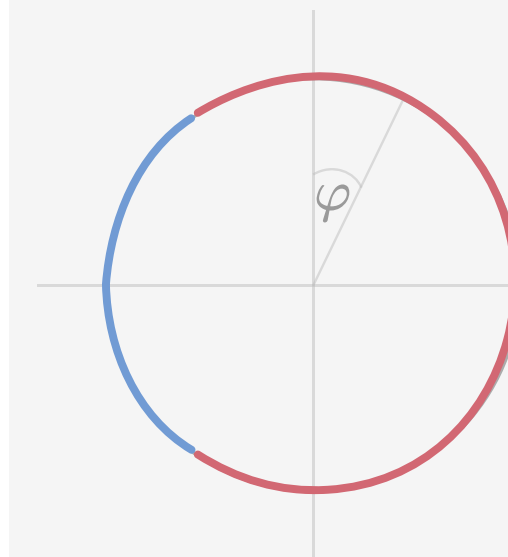
$$\mathcal{D} = 0 \Leftrightarrow \text{finitely many } \blacksquare$$

Theorem 2. \mathcal{D} can be computed to arbitrary precision.

Theorem 3. “ $\mathcal{D} \in \mathbb{Q}$?” is decidable, when there are at most three dominant eigenvalues.

III. The Example, or First Observation

$$\varphi = \cos^{-1} 4/5$$



How frequently are we on the red arc?

Theorem (Weyl, 1910). Let ρ be an irrational real number. Then the sequence: $\rho, 2\rho, 3\rho, \dots$ is uniformly distributed mod 1.

$$\mathcal{D} = \frac{\text{length of red arc}}{2\pi} = \frac{\cos^{-1}(-2/3)}{\pi} = 0.732278\dots$$

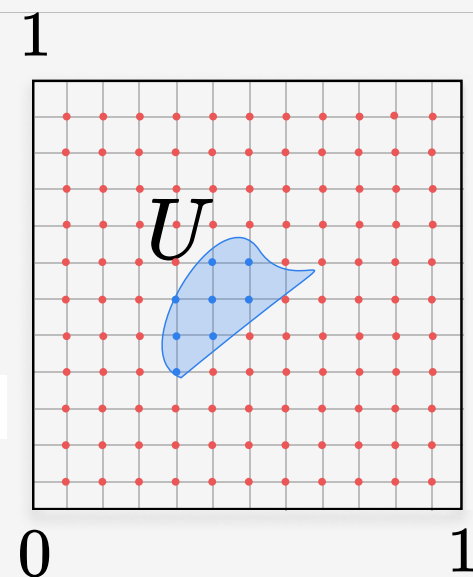
IV. The Proof

Approximating the Density

1. How to make a grid such that $\bullet \in U$ can be decided

2. How fine should the grid be for $|\text{approx - vol}| < \epsilon$?

Pascal Koiran:
Approximating the Volume of Definable Sets. FOCS 1995: 134-141



Approximate volume \approx $\frac{\text{number of red dots}}{\text{number of blue dots}}$

Theorem 2. \mathcal{D} can be computed to arbitrary precision.

V. The Open Problem

Decide whether $\mathcal{D} \in \mathbb{Q}$

When there are no multiplicative relations:

Decide if $\int_{\mathcal{L}} \prod_{i=1}^{\eta} \frac{1}{\sqrt{1-x_i^2}} d\vec{x} \in \mathbb{Q} \pi^{\eta}$

polytope